

A STUDENT INFORMATION SYSTEM

FOR

MICROCOMPUTERS

A THESIS

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This study is dedicated to my daughter Trisha. May she accept and conquer the challenges of the future.

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CHAPTER I

INTRODUCTION

Data base management systems (DBMS), and management information systems (MIS) are terms that are frequently used in modern computer installations. Although both concepts originated with the automation of accounting, production, and inventory control, they evolved into major systems as the emphasis in computer processing shifted in an attempt to meet management's information needs.

The difference between data base management systems and management information systems is subtle but real. C. J. Date (1) defines a data base management system as a collection of stored operational data used by the application systems of some particular enterprise. An appropriate definition of an management information system is one presented by Riley (2). An information management system is a system for collecting, sorting, retrieving, and processing information which is used, or desired, by one or more managers in the performance of their duties. Thus a data base management system is part of a larger information management system. This study will focus on data base management systems and will investigate the feasibility of

using microcomputers to implement and maintain a small, stand-alone data base system.

George M. Scott (3) provides a definition of a modern data base as a collection of computer files of data structured to enable efficient updating, maintenance, reporting, and storage of data and to enable rapid retrieval of all stored data that must be brought together for a particular operation or managerial purpose. This definition implies two major dimensions of a data base. The first dimension relates to structuring data files to enable efficient updating, maintenance, reporting, and storage of data. Data base structuring involves complex data-file design and data retrieval technology. The various data structural models discussed in this study include the hierarchical or tree model, the network model, and the relational model.

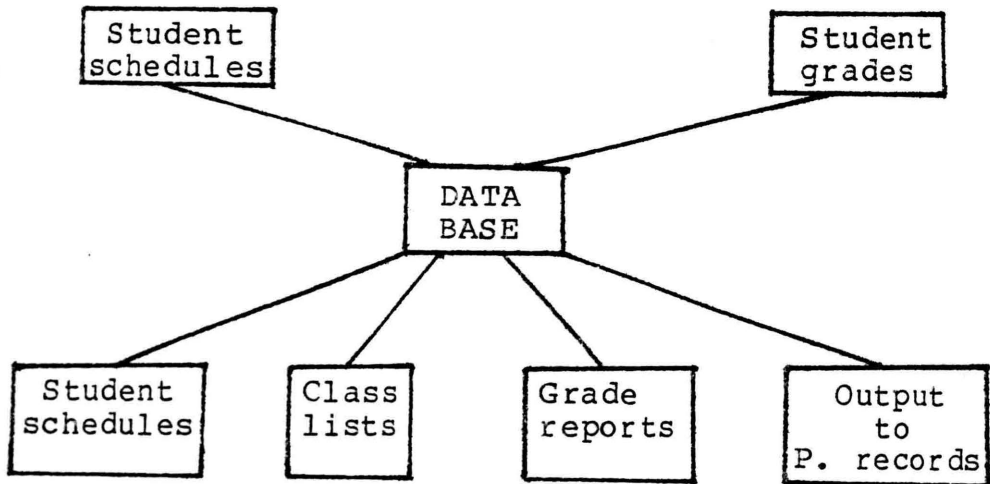
The second dimension of data bases referred to in Scott's definition, enabling rapid retrieval of all stored data that must be brought together for a particular operation or managerial purpose, illustrates the relationship that exists between data base systems and management information systems. The functions of query languages and the role they play in data base management design is also included within this second dimension. This

study is concerned with the design and implementation of a data base system. The theoretical discussion of query languages, their implementation, and completeness is not included.

To design any data base management system, it is necessary to define the information to be represented within the system. A. T. F. Hutt (4) defines a "Perceived Field" of a data base as that part of the real world which is to be reflected within the system. Hutt stresses the importance of deciding on the size of the perceived field and the necessity of allowing the field of perception to grow via a number of phased expansions of the data system. The "Field of Perception" of the data base to be implemented in this study is outlined in figure 1. Since the purpose of the study is to investigate the feasibility of using micro computers to implement and manage a stand-alone data base, a small environment has been selected. The data base to be implemented is to consist of the schedules of students, class lists, and student grade reports for a high school environment of five hundred students.

In designing a data base management system, several basic functions of the system must be considered. The system must be able to retrieve particular records, insert new records, delete individual records, and update specific

Input data



Information retrieved from the system

Procedures:

1. Input student schedules.
2. Produce class lists.
3. Input student grades. (mid and end of semester)
4. Produce student grades reports.
5. Prepare output for the permanent record file.

note: An environment of five hundred students is assumed and although this implementation does allow for reporting to the permanent record file, this file is not included in the system system as developed in this study.

Figure 1: Perceived Field

information represented within the system. Fundamental to the design of the data manipulation language (DML) which supports these basic functions is the conceptual view of the data. In particular, the data structure dictates the design of the corresponding data manipulation language.

At present, the three best known data structural models are: the hierarchic or tree model, the network model, and the relational model. The conceptual view of data suggested by these models dictates the design of the DML needed to support the basic functions of inserting records into the system, deleting data from the system, and updating information contained within the system. The hierarchic or tree model employs a hierarchical design of data and was adapted by the Conference of Data Systems Languages (CODASYL) and presented in their 1971 DBTG report. This model suggests that real world data lies within hierarchical structures with one record at the top of the tree, referred to as the root node, and other records dependent on it. In general any record may have any number of dependents and each of these dependents may have any number of dependents, and so on. In the hierarchical model each dependent must have exactly one successor. Figure 2 illustrates a hierarchical view of the data represented in a student information system.

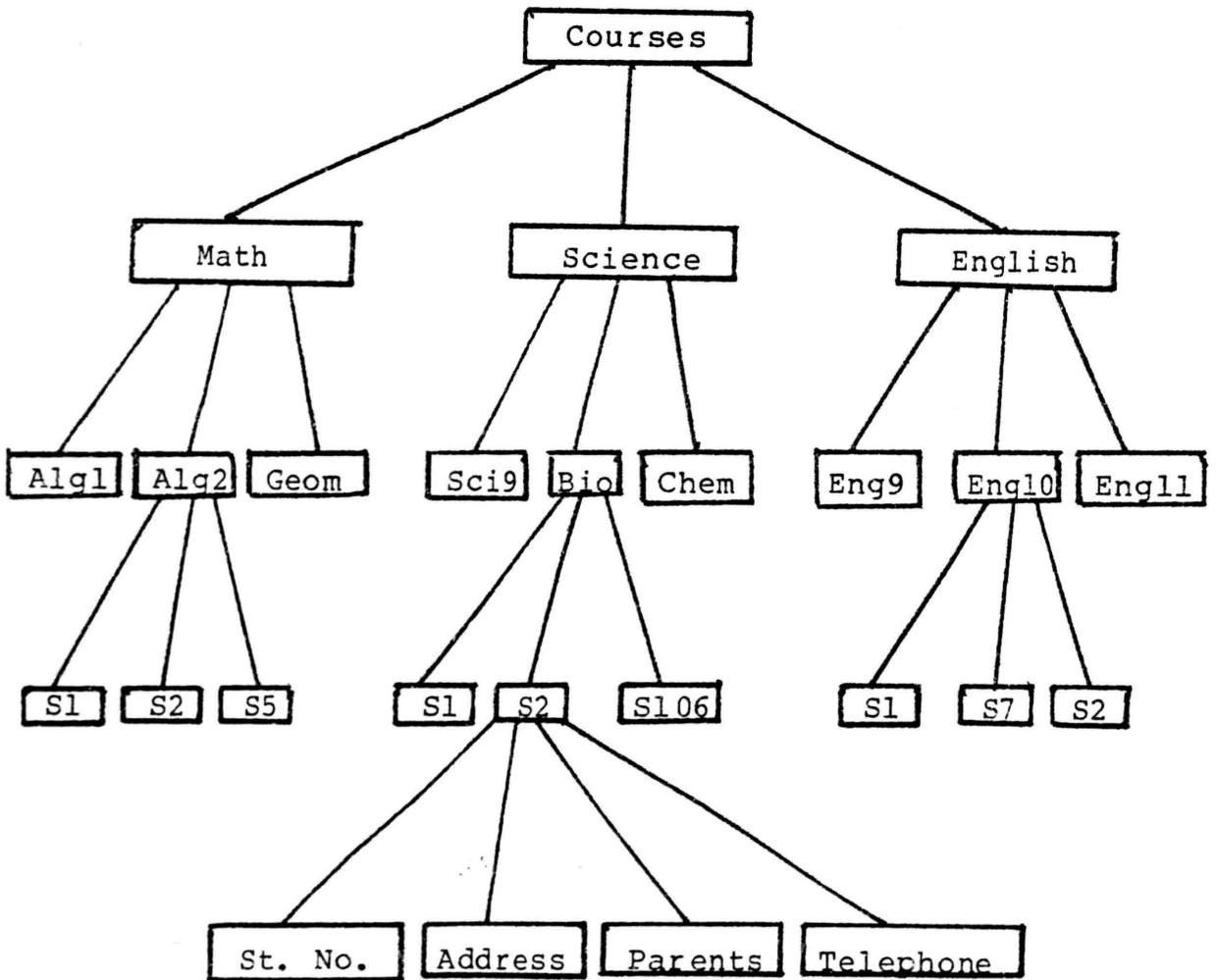


Figure 2: Hierarchical View of Student Information Data

✓ The DML needed to maintain the basic functions of insertion, deletion, and updating of information becomes complex. It is impossible to insert a new student until his classes have been defined. The procedure required to delete an existing student is also complex since the data base must be searched in its entirety and all occurrences of the particular student must be deleted. It also follows that if we delete all classes of a particular student, we have deleted the existence of the student. Thus the insert and delete anomalies reflect a similar problem. The update procedure also involves searching the entire structure. To update a particular student's address, for example, the entire structure must be searched and each occurrence of the address modified.

✓ The network data structural model is a more general structure than the hierarchical model. A given record in the network model may have any number of immediate superiors whereas a record is restricted to exactly one superior in the hierarchical model. The network approach allows many-to-many relationships whereas the hierarchical model is restricted to one-to-many relationships. Figure 3 illustrates a network view of the student record data representation. The procedures for inserting, deleting, and updating information become conceptually simple within the

network model and the anomalies discussed with the hierarchical model do not arise. A new student could be added to the structure before his schedule is defined and to delete a student, or update a student's address we need not search the entire structure. The prime disadvantage of the network model, however, lies in the complexity of the DML required. The programming involved becomes extremely sophisticated resulting in relatively high implementation and maintenance costs.

✓ The relational approach, first proposed by Codd (5) in 1971, is a different approach to describing and manipulating data. It views the data base as a simple collection of two dimensional tables called relations. Although conceptually simple, as illustrated in figure 4, a relation can be explicitly defined as a two dimensional table with m rows, each made up of a set of n -tuples. Each of the columns in a relation is a set of values of one attribute and is referred to as a domain. The power of a relational data base exists in the data manipulation language rather than in the data structure. Figure 4 illustrates how the data in the student information system can be represented by a number of relations. The inserting, deleting, and updating procedures become relatively simple. A new student can be inserted into the student-schedule or student-information relations

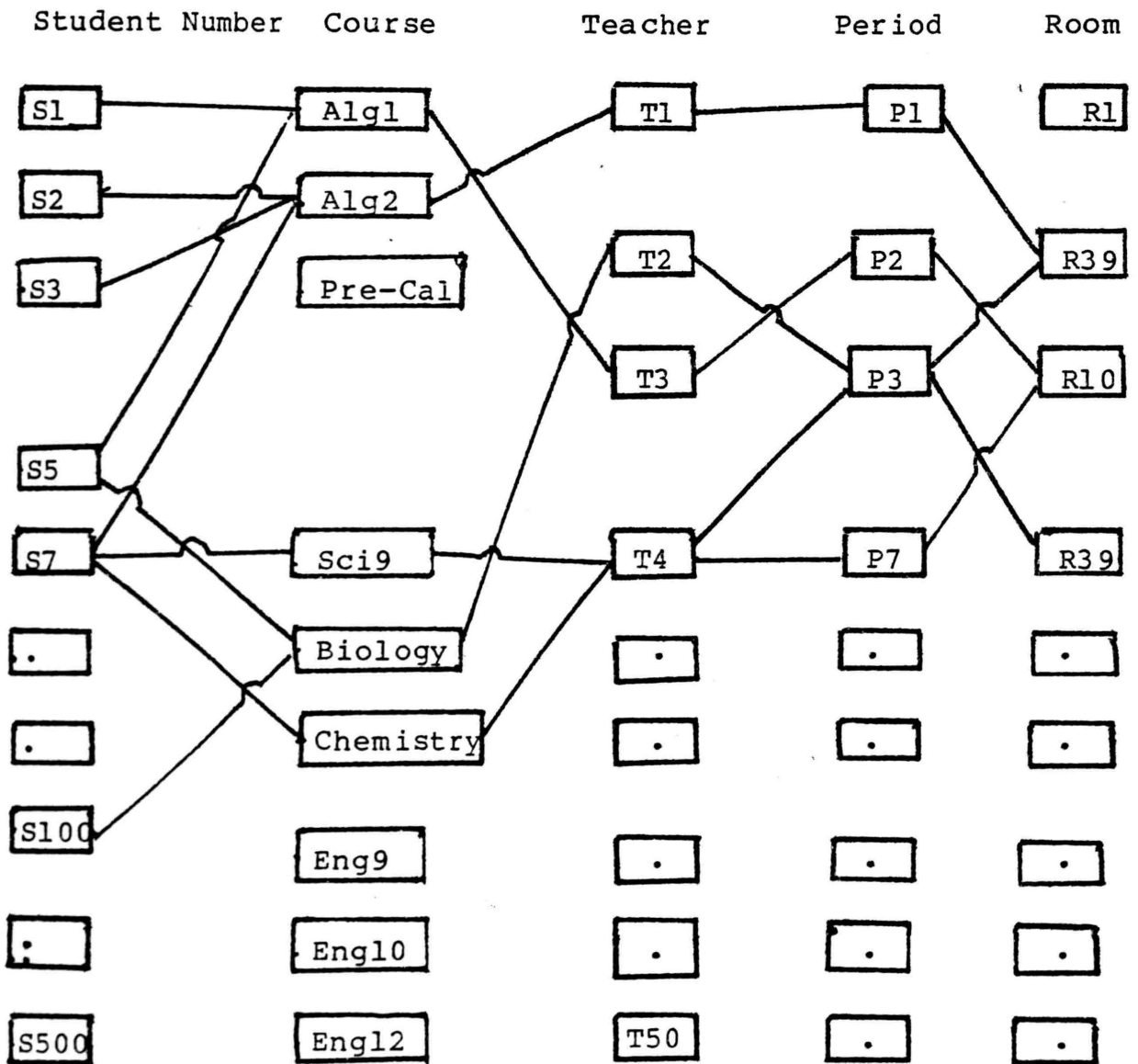


Figure 3: Network View of Student Information Data

Student-Schedule Relation:

Student Number	courses			
S1	Alg1	Engl0	Biology	P.E.
S2	Eng9	French	Sci9	St. Hall
S3	Engl2	Physics	Pre-Cal	German
S4	Hist9	Alg1	Biology	French
S5	Am.Prob.	P.E.	Engl0	Geometry
.
.
S500

Student-Information Relation:

Student Number	Address	Parents	Telephone
S1	521 Windsor Dr.	John Jones	383-3995
S2	436 Johnson Dr.	Jane Smith	452-6547
.	.	.	.
.	.	.	.

Figure 4: Relational View of Student Information Data

as required, a student can be deleted from any relation without affecting other relations, and procedures to update information require single operations. ✓

X A historical look at the development of Data Base Management Systems shows that although recent trends appear to be moving toward the relational view of data, commercially available systems generally employ the hierarchical or network data structural model. Figure 5, taken from Data Base Management Systems (6), outlines the general data bases that were commercially available in 1977 and illustrates that the CODASYL hierarchical or tree structural design is the most widely used data model. IBM's Information Management System (IMS2) also uses a hierarchical data model but does not follow CODASYL standards and thus receives a classification of its own. Although there are no general relational data base systems commercially available, several experimental relational systems have been implemented with promising results. IBM's IS/1 shows the feasibility of supporting relational algebra, a data manipulation language proposed by Codd (5), and the experimental SYSTEM R, being developed at the IBM Research Laboratory, does classify as a pure relational system.

The specific aims of this study are to investigate the feasibility of using microcomputers to implement and

SYSTEM	SUPPLIER	CLASSIFICATION
DMS-11	Burroughs	Network
DMS-170	Control Data	CODASYL DBTG
TOTAL	CINCOM	Network
IDMS	Cullinane	CODASYL DBTG
DBMS/10	DEC	CODASYL DBTG
IMS2	IBM	Hierarchical
SYSTEM 2000	MRI	Tree, Inverted
ADABAS	Software AC	Network inverted
DMS 1100	UNIVAC	CODASYL DBTG

Figure 5: A partial list of Data Management Systems commercially available

maintain a stand-alone data base to be used in a high school environment. A data base management system, with the field of perception as outlined in figure 1, will be implemented using the relational data model. The multi-step design methodology, presented by Hutt (4), is followed. The steps

involved in Hutt's multi-step methodology are:

1. THE OBJECT ENVIRONMENT: The perceived field is identified and real world objects that exist within it are defined.
2. THE INFORMATION ENVIRONMENT: A number of relations in third normal form, as defined in chapter two of this document, are defined.
3. THE STORED ENVIRONMENT: The indexes and storage structures are defined.
4. THE SOFTWARE ENVIRONMENT: The development of the computer software which operates on the stored data to perform the required results.

CHAPTER II

THE ARCHITECTURE

2.1 Definitions

C. J. Date (1) defines a relation R on a collection of N sets D_1, D_2, \dots, D_n as a set of ordered n -tuples $\langle d_1, d_2, \dots, d_n \rangle$ such that d_1 belongs to D_1 , d_2 belongs to D_2, \dots, d_n belongs to D_n . Sets D_1, D_2, \dots, D_n are called the domains of the relation R and the value of n is the degree of the relation. It is important to distinguish between a domain and an attribute which is drawn from the domain. An attribute represents the use of a domain within a given relation. To emphasize this distinction; the attributes for the relation Student-Schedule given in figure 4 of chapter 1 consist of course1, course2, course3, and course4, which are all taken from the domain of courses offered. Generally speaking, if a relation is considered to be a table of values with m rows and n columns, the attributes can be considered to be the names given to the individual columns.

In any relation there will be at least one attribute, or a combination of attributes, that are unique within the relation. If the relation contains one such attribute, this attribute is known as the primary key of the relation. In

the relation Student-Schedule the attribute (student number) serves as the primary key since each occurrence of a specific student number uniquely defines a 5-tuple. If a relation contains more than one attribute combination possessing this unique identification property, one attribute combination is arbitrarily chosen as the primary key while the remaining "candidate keys" are referred to as alternate keys.

The problem of defining what relations are needed and what their attributes should be is the fundamental problem of designing a relational data base. The concept of normalization of relations plays an important role in defining appropriate relations to adequately represent the data. The following discussion introduces the concept of normalization and discusses why it is important to properly define the relations in third normal form when designing a relational data base.

Every relation is considered to be in first normal form if each entry in the relation is atomic (i.e. if each entry is nondecomposable as far as the system is concerned). This definition is fundamental in defining relations and simply means that each entry in the table consists of precisely one value, never a set of values. A relation that is only in first normal form leads to anomalies similar to those

encountered by hierarchies discussed in chapter 1. For example, the relation given in figure 6(a) is in first normal form and problems occur with the three basic operations of inserting, deleting, and updating information. We cannot enter the fact that a new student exists until he has defined at least one class, if we delete all classes we delete the student, and if the teacher of a particular course is changed the entire structure must be searched and all appropriate changes made.

In order to resolve some of these difficulties, first normal form relations are transformed into second normal form relations. Before we can define second normal form relations, the concepts of functionally and fully functionally dependencies must be understood. An attribute Y of a relation R is said to be functionally dependent on attribute X of R if and only if each X -value in R is associated with precisely one Y -value in R . Furthermore, attribute Y is fully functionally dependent on attribute X if it is functionally dependent on X and not functionally dependent on any proper subset of X . In the relation given in figure 6(a) the attribute (Sname) is functionally dependent on the primary key (student number, course number) but it is not fully functionally dependent on this attribute since (Sname) is also dependent on (student number) which is

a proper subset of the attribute combination (student number, course number).

We now define a relation to be in second normal form if and only if it is in first normal form and every nonkey attribute is fully functionally dependent on the primary key. It is clear that the relation Student-Information given in figure 6(a) is not in second normal form since the nonkey attribute (Sname) is not fully dependent on the key (Snumber, Cnumber). A relation in first normal form can always be transformed into equivalent relations in second normal form. The process involves eliminating the non-full functional dependencies. The functional diagram given in figure 6(b), gives functional dependencies and leads to defining the relations in second normal form given in figure 7. The anomalies discussed above do not exist with relations in second normal form since a student can exist before classes are defined by simply inserting the appropriate tuple into the student-name relation. The relation Teacher(Cnumber, Tnumber, Tname), however, still creates problems. To be specific, the dependency of (Tname) on the primary key is transitive. Each (Cnumber) determines a (Tnumber) which in turn determines a (Tname). The anomaly created by the transitivity exists in that we cannot insert a teacher into the data base until a specific

class has been defined for him to teach.

To overcome the above anomaly we further reduce the second normal form relations to third normal form. A relation is in third normal form if and only if it is in second normal form and every nonkey attribute is nontransitively dependent on the primary key. This reduction process is accomplished by eliminating transitive dependence and the relation `Teacher(Cnumber,Tnumber,Tname)` is reduced to two relations. The relations `CT(Cnumber,Tnumber)` with primary key (Cnumber) and `TN(Tnumber,Tname)` with primary key (Tnumber) replace the relation Student-Information. Figure 8 gives the four relations in third normal form equivalent to the given Student-Information relation.

Several other normal forms have been defined which deal with relations with more than one candidate key but since all relations used in the design and implementation of the data base discussed in this study contain a single key, relations defined in third normal form will ensure that the above anomalies will not occur when implementing the insert, delete, and update procedures.

2.2 THE OBJECT ENVIRONMENT

As illustrated in figure 1, the student courses and

grades are to be input into the system and the system is to produce student schedules, class lists, semester grade reports, and output to the permanent record file. This system does not include the permanent record file but produces output containing class codes, grades, teacher codes, and teacher comment codes. The object environment of the relational data base to be implemented is:

Student number: A five digit integer. The first two digits identify the graduation year and the remaining three digits identify the student. Student number 85123, for example, identifies that student number 123 will graduate in 1985.

Course number: A three digit integer representing a particular course. The specific departments are identified by the leading digit and individual courses within the department are identified by the remaining two digits. In this study the departments are identified by the following scheme.

- 1 identifies an elective.
- 2 identifies an English course.
- 3 identifies a history course.

- 4 identifies a mathematics course.
- 5 identifies a physical education course.
- 6 identifies a science course.

This identification is random and has no bearing on the design of the system. A course such as 421 will imply mathematics course number 2, section 1.

Teacher Number: A two digit integer identifying the teacher of a particular course.

Grades: An integer, maximum of two digits, representing the student grade in a particular course. The system allows for a mid-semester and end of semester grades.

Teacher comment: A one digit integer (0-9) representing a teacher comment code. Two comments are represented; a mid-semester comment and a end of semester comment.

Room Number: A two digit number identifying the room number for a particular class.

2.3 THE INFORMATION ENVIRONMENT

The information is represented within the system by a number of relations in third normal form. The relations

required are the Student-Schedule relation, the Teacher-Schedule relation, and the Class-Information relation. Although the relation Student-Name is given here for completeness it is not included within the system as implemented in this study.

The Student-Schedule relation has domains student number, course number, grades, and teacher comment. Since a seven period day is assumed, the domains course number, grades, and teacher comment are repeated to form several attributes. Figures 9 and 10 give a detailed description of these relations including the primary keys and descriptions of the attributes of each of the relations.

2.4 THE STORED ENVIRONMENT

In this section the storage structures are defined. The Student-Schedule relation is stored externally as a random access file with a defined record length of one hundred and four bytes. Each record contains one 36-tuple of the relation as outlined in figure 11. Since the system is designed to accommodate a maximum of five hundred students, this file contains five hundred records. The file is indexed by the key attribute student-number. The index is structured as a binary tree and is stored in core memory as a five hundred by four array and on auxiliary disk as a

sequential file.

The Teacher-Schedule relation has a structure similar to the Student-Schedule relation with each 8-tuple stored as one record of a random access file. This file has a defined record length of thirty one bytes and figure 12 gives a detailed description of one record of this file. The Teacher-Schedule index is also structured as a binary tree and is stored in core memory as a fifty by four array and on the disk storage as a sequential file.

The Class-Information relation with attributes (Cnumber, Tnumber, Rnumber, Csize, S1, S2, ..., S35) is also stored as a random access file where each record has a defined length of two hundred and twenty five bytes. Each record of the file stores the course number, teacher number, period number, room number, class size, and the list of students currently enrolled in the class. The Class-Information index is also structured as a binary tree and is stored in core in an one hundred by fifty array and on the disk as a sequential file.

-a-

Relation: Student-Information

Attributes (Snumber, Sname, Cnumber, Tnumber, Tname)

Primary key (Snumber, Cnumber)

Snumber	Sname	Cnumber	Tnumber	Tname
S1	John Smith	C7	T6	Mr. White
S5	Jack Jones	C9	T1	Ms. Brent
S46	Sam Bright	C7	T6	Mr. White
.
S498	Howard Go	C18	T15	Mr. Black

-b-

Functional Dependencies

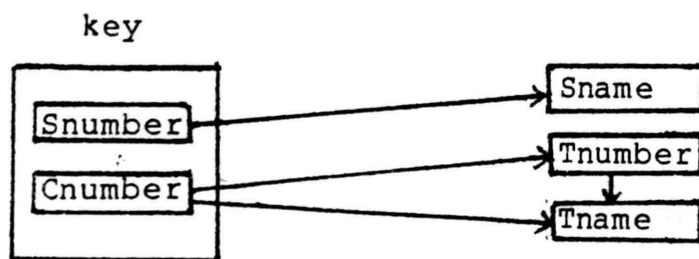


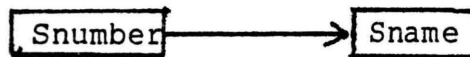
Figure 6: Relation Student-Information in
"First Normal Form"

-a-

Relation: Student-Name
 Attributes: (Snumber, Sname)
 Key: (Snumber)

Snumber	Sname
S1	John Smith
S5	Jack Jones
S45	Sam Bright
S498	Howard Go

Function
 Dependency



-b-

Relation: Teacher
 Attributes: (Cnumber, Tnumber, Tname)

Cnumber	Tnumber	Tname
C7	T6	Mr. White
C9	T1	Ms. Brent
C7	T6	Mr. White
C18	T15	Mr. Black

Functional Dependencies

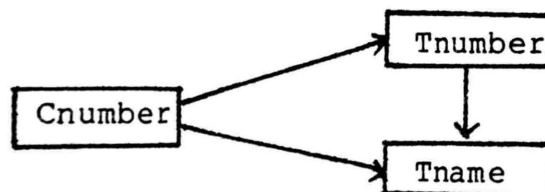


Figure 7: Relation Student reduced to "Second Normal Form"

Relation: Student
 Attributes: (Snumber, Sname)
 Key: (Snumber)

Snumber	Sname
S1	John Smith
S5	Jack Jones
.	.
S498	Howard Go

Relation: Schedule
 Attributes: (Snumber, Cnumber)
 Key: (Snumber)

Snumber	Cnumber
S1	C7
S5	C9
.	.
S498	C18

Relation: CT
 Attributes: (Cnumber, Tnumber)
 Key: (Cnumber)

Cnumber	Tnumber
C7	T6
C9	T1
.	.
C18	T15

Relation: TN
 Attributes: (Tnumber, Tname)
 Key: (Tnumber)

Tnumber	Tname
T6	Mr. White
T1	Ms. Brent
.	.
T15	Mr. Black

Figure 8: Relations in "Third Normal Form"

-a-

Relation: Student-Schedule

Attributes:

..... period 1 period 2
 (SNumber, cl, g1l, coll, glf, colf, c2, g2l, co2l, g2f, co2f,
 period 7
 ,g7, g7l, co7l, g7f, co7f)

Key: Snumber

ci....implies...course number corresponding to
 appropriate period (7 period day).

gil....implies...the mid-semester grade.

coil..implies...the mid-semester comment.

gif...implies...the end of semester grade.

coif..implies...the end of semester comment.

-b-

Relation: Teacher-Schedule

Attributes: (Tnumber, C1, C2, C3, C4, C5, C6, C7)

Key: Tnumber

Tnumber...implies...Teacher number

Ciimplies...Course number of the ith class
 period.

Figure 9: Relations Student-Schedule and Teacher-Schedule

-a-

Relation: Class-Information

Attributes: (Cnumber, Tnumber, Pnumber, Rnumber, Csize, S1, ..., Sn)

Key: Cnumber

Cnumber...represents the course number.

Tnumber...represents the teacher number.

Pnumber...represents the period number.

Rnumber...represents the room number.

Csize....represents the current class size.

Si.....represents a student number.emacs

-b-

Relation: Student-Name

Attributes: (Snumber, Sname, Pname, Address, Telephone)

Key: Snumber

Snumber...represent the student number.

Sname.....represents the student name.

Pname.....represents the parent's name.

Address...represents the student's address.

Telephone.represents the student's phone number.

Figure 10: Relations Class and Student


```

.5.4.) .9.9.9.) .5.1.1.) .5.1.2.) .9.9.9.5.1.3.) .5.1.4.) .5.1.
0      5      10     15     20     25
.1.) :
31

```

bytes: 0,1	teacher number.
2	end of field marker.
3,4,5	course (period 1).
6	end of field marker.
.	.
.	.
30	end of record marker.

Figure 12: BYTE MAP OF TEACHER SCHEDULE FILE

```

.5.1.1.) .5.4.) .2.) .2.2.) .1.8.) : : : : : : : : :
0      5      10     15     20     25

: : : : : : : : : : : : : : :
30      35      225

```

bytes:	0,1,2	course number.
	3	end of field marker.
	4,5	teacher number.
	6	end of field marker.
	7	period number.
	8	end of field marker.
	9,10	room number.
	11	end of field marker.
	12,13	class size.
	14	end of field marker.
	15-224	student numbers (maximum of 35).
	225	end of record marker.

Figure 13: BYTE MAP OF CLASS-INFORMATION FILE

CHAPTER III

THE SOFTWARE ENVIRONMENT

The student information system developed in this study consists of a series of eight programs as outlined in figure 14 and listed in appendices A through H. Each program is structurally designed and documentation on the programs and subroutines is given in this chapter. Although most of the procedures are self explanatory, two procedures are documented separately in sections 3.1 and 3.2 of this chapter.

3.1 THE INDEX STRUCTURES

The system programs access the random access files by storing the key elements and file record numbers in an index structured as a binary tree. The index structures are stored externally as sequential files and internally as n by four arrays. The four entries per element represent the left pointer, key element, right pointer, and the random access file record number corresponding to the key element. Figure 15 gives the internal matrix representation of a random selection of class numbers in terms of both the matrix and the abstract binary tree representation. It is

possible to search the index to find a given key by following pointers in the array structure. Figure 16 gives the BASIC program for searching the matrix representation of a binary tree.

3.2 THE BALANCE PROCEDURE

The advantage of using a binary tree structure to store an index is a reduction in time required to search the index. The time required to search a binary tree is directly proportional to the length of the tree. Since it is possible to construct various binary trees using the same nodes, as illustrated by figure 17, it is important to optimize the structure of the tree by minimizing its length. The student information system developed in this study contains a tree balancing procedure which produces the optimal tree structure, a balanced binary tree. The procedure, included within the BUILD program listed in appendix III, utilizes the in-order, length, and position pointers subroutines and the algorithm for the procedure is documented below. Figure 18 further documents the balance algorithm by presenting the step by step calculations performed when the index given in figure 17(a) is balanced to produce the structure presented in figure 17(b).

STEP 1: Traverse the tree structure in left-node-right order

STEP 1: Traverse the tree structure in left-node-right order and store the row numbers in a vector ODR.

STEP 2: Calculate the maximum length of the largest full balanced tree that can be created with the given nodes. The maximum full balanced tree size, number of elements, will be a positive integral power of two minus one.

STEP 3: Determine the number of elements to temporarily delete from the ODR vector, store the first odd elements in a vector OMIT and pack ODR.

STEP 4: Create a full balanced tree by inserting appropriate pointers into the structure.

STEP 4a: Set the left and right pointers of all odd numbered nodes to zero. That is nodes $ODR(1), ODR(3), \dots, ODR(2i-1)$ left and right pointers are set to zero.

STEP 4b: Place pointers in the index matrix in a number of passes as outlined below.

PASS 1.-n = 2.

-left pointer of node $ODR(n)$ = node $ODR(n-1)$.

-right pointer of node $ODR(n)$ = node $ODR(n+1)$.

-n = n + 4.

-repeat until n is greater than the length

of packed ODR.

PASS 2.-n = 4.

- left pointer of node ODR(n) = node ODR(n-2).
- right pointer of node ODR(n) = node ODR(n+2).
- n = n + 8.
- repeat until n is greater than the length of
of packed ODR.

PASS I: where 2 to the power I is less than the length
of packed ODR.

- n = 2 to the power I.
- K = n + 1.
- left pointer of node ODR(n) = node $n - 2^{(i-1)}$.
- right pointer of node ODR(n) = node $n + 2^{(i+1)}$.
- n = n + 2 to the power K.
- repeat until n is greater than the length of
packed ODR.

STEP 5: Place the omitted nodes into the tree by extending
the lower left nodes.

- left pointer of node ODR(1) = node omit(1)
- right pointer of node ODR(1) = node OMIT(2).
- left pointer of node ODR(2) = node OMIT(3).
- continue until all nodes deleted in step 3 are
inserted.

STEP 6: Set both left and right pointers of the inserted nodes to zero.

STEP 7: Define the root node. The root node will be node $ODR(k)$, where $k = (\text{length of packed ODR} + 1)$ divided by two.

3.3 THE SYSTEM PROGRAMS

The eight programs of the student information system are documented in this section with all subroutines documented in section 3.4. Figure 14 outlines the information system in terms of the programs accessed by the system and figure 19 gives a description of the fundamental variables used in the system software.

3.3.1 PROGRAM SYSTEM MENU: This program, listed in appendix A, controls the system software. The menu, which consists of the system programs, appears on the screen when the system is booted and the user selects which program to execute.

3.3.2 PROGRAM INITIALIZE FILES: This program, listed in appendix B, initializes the random access files as outlined in figures 11, 12, and 13. All fields are set equal to zero to allow for error processing. The index files are

opened and the first three fields of each index file are set to zero. These fields contain the root node of the binary index, the number of elements in the index, and the number of elements in the index contained when it was last balanced.

3.3.3 PROGRAM BUILD SCHEDULES: This program, listed in appendix C, builds the student and teacher schedule files. The program accepts the schedules from the user, enters the key element into the appropriate index, and writes the schedule to the correct file and record. The class index is also maintained by this program and the class information file is updated as a given teacher or student is added to the data base. Figure 20 gives the structural design of the build schedules program.

3.3.4 PROGRAM RETRIEVE SCHEDULES: This program, outlined in figure 20 and listed in appendix D, allows the user to print one student's schedule, one teacher's schedule, all student schedules, or all teacher schedules. If the user chooses to print all schedules the

schedules will be produced in order of increasing student or teacher number.

3.3.5 PROGRAM UPDATE: This program allows the user to make changes in the student schedules, student's grades and teacher comments, teacher schedules, or room numbers. The design, by subroutines, is outlined in figure 22 and the program listing is given in appendix E.

3.3.6 PROGRAM ENTER GRADES: This program asks the user if end of semester or mid-semester grades and comments are to be entered, produces class lists, inputs the grades and comments, and writes the student grades and comments to the student schedule file. Figure 23 gives the structural design of this program and a program listing of the enter grades program is given in appendix F.

3.3.7 PROGRAM PRINT GRADES: A program which prints either the mid-semester or the end of semester grade reports. The program will produce grade reports for an individual student, or all grade reports will be produced in order of increasing student number. The structural design of the print grades

program is given in figure 24 and a listing of the program is included in appendix G.

3.3.8 PROGRAM RETRIEVE CLASS LISTS: This program allows the user to retrieve class lists. The user can access lists for one class, all classes for one teacher, or all classes that exist within the data base. Figure 25 gives the structural design of this program and the program listing is given in appendix H.

3.4 SUBROUTINE DOCUMENTATION

The various subroutines accessed by the programs of the student information system are documented in terms of the function performed by each subroutine. The subroutines are standard in design and the documentation given presents an overview of the purpose of the subroutine rather than a detailed description of the algorithms employed.

3.4.1 SUBROUTINE READ TEACHER INDEX: A subroutine which reads the teacher index from the sequential teacher index file and internally stores the index as the TINDX array.

- 3.4.2 SUBROUTINE READ STUDENT INDEX: A subroutine which reads the student index from the sequential student index file and stores the index internally as the SNDX array.
- 3.4.3 SUBROUTINE READ CLASS INDEX: A subroutine which reads the class index from the sequential class index file and internally stores the index in the CINDX array.
- 3.4.4 SUBROUTINE CHANGE SCHEDULE: This subroutine allows the user to make changes in the schedule during the input process. This subroutine calls the display schedule, change number, and change classes subroutines.
- 3.4.5 SUBROUTINE DISPLAY SCHEDULE: This subroutine is called by the change schedule subroutine and displays the current teacher or student schedule.
- 3.4.6 SUBROUTINE CHANGE NUMBER: This subroutine allows the user to correct the teacher or student number before the number is stored in the appropriate file.
- 3.4.7 SUBROUTINE CHANGE CLASSES: This subroutine allows the user to change the classes, by period, in the schedule before it is stored in the

appropriate file.

- 3.4.8 SUBROUTINE FIND STUDENT RECORD: This subroutine searches the student index and inserts the given student into the student index structure. The corresponding record number of the student schedule file is also recorded in the index.
- 3.4.9 SUBROUTINE WRITE STUDENT SCHEDULE: A subroutine that writes the given student's schedule to the appropriate record of the student schedule file.
- 3.4.10 SUBROUTINE UPDATE CLASS INDEX: This subroutine searches the class index and returns the corresponding record number of the class information file if a given course exists in the index. If the course does not exist the class is added to the index and the course and period numbers are printed to the appropriate record of the class information file.
- 3.4.11 SUBROUTINE REVISE CLASS INFO (STUDENT): A subroutine which increments the class size and adds the given student to the class list.
- 3.4.12 SUBROUTINE REVISE CLASS INFO (TEACHER): A subroutine

which adds the given teacher to the appropriate record of the class information file.

3.4.13 SUBROUTINE EXIT: This subroutine controls the exit process. All index structures are balanced if more than five elements have been added to the index since it was last balanced.

3.4.14 SUBROUTINE BALANCE INDEX: This subroutine controls the balance procedure as outlined in section 3.2 of this chapter by calling the in order, length, and the position pointers subroutines.

3.4.15 SUBROUTINE IN-ORDER: A subroutine to traverse the binary tree index structure in (left-node-right) order and produce the vector ODR. The vector ODR contains pointers to the node elements of the index tree to allow processing in order of increasing student, teacher, or course numbers.

3.4.16 SUBROUTINE POSITION POINTERS: This subroutine uses the ODR vector formulated by the in-order subroutine to build a balanced tree structure. This subroutine uses the

algorithm presented in section 3.2.

3.4.17 SUBROUTINE LENGTH: This subroutine calculates the size of the largest balanced tree that can be constructed.

3.4.18 SUBROUTINE WRITE TEACHER INDEX: This subroutine writes the teacher index to the teacher index file.

3.4.19 SUBROUTINE WRITE CLASS INDEX: This subroutine writes the teacher index to the class index file.

3.4.20 SUBROUTINE WRITE STUDENT INDEX: This subroutine writes the student index to the student index file.

3.4.21 SUBROUTINE RETRIEVE STUDENT SCHEDULE (ONE STUDENT):
This subroutine reads the student schedule for a given student from the student schedule file.

3.4.22 SUBROUTINE FIND CLASS RECORD: A subroutine which searches the class index and returns the corresponding record number of the class information file.

3.4.23 SUBROUTINE READ TEACHER-ROOM NUMBERS: A subroutine which reads the teacher and room number from the appropriate record of the class information file.

3.4.24 SUBROUTINE PRINT STUDENT SCHEDULE: A subroutine to display the student schedule for a given student. The teacher and room numbers are also displayed.

3.4.25 SUBROUTINE RETRIEVE TEACHER SCHEDULE (ONE TEACHER): This subroutine reads the teacher schedule for a given teacher from the teacher schedule file.

3.4.26 SUBROUTINE RETRIEVE STUDENT SCHEDULES: A subroutine to control the processing of all student schedules. The subroutine first calls subroutine in order and then processes each student's schedule by reading the schedule and calling the find class record, read teacher-room numbers, and print student schedule subroutines.

3.4.27 SUBROUTINE RETRIEVE TEACHER SCHEDULES: A subroutine to control the processing of all teacher schedules. The subroutine calls subroutine in-order to produce a list of all teachers in order of increasing teacher number. The class schedule for each teacher is then processed by calling subroutines read teacher schedule, read teacher-room numbers,

and print teacher schedule.

3.4.28 SUBROUTINE PUSH STACK: This subroutine is used by subroutine in-order. The current node and a direction (D = 1 represents a left pointer and D = 2 represents a right pointer) are pushed onto the stack.

3.4.29 SUBROUTINE POP STACK: This subroutine returns the node and direction from the top of the stack.

3.4.30 SUBROUTINE PRINT TEACHER SCHEDULE: A subroutine to print the teacher schedule and room numbers for a given teacher.

3.4.31 SUBROUTINE BYTE POINTERS MID-SEMESTER: A subroutine to set B1 and B2 to ten and thirteen respectively to allow access to the mid-semester grades and comments.

3.4.32 SUBROUTINE BYTE POINTERS FINIAL: A subroutine to set B1 and B2 to fifteen and eighteen respectively to allow access to the end of semester grades and comments.

3.4.33 SUBROUTINE FIND TEACHER RECORD: A subroutine to search the teacher index for a given teacher and return the record number of the teacher schedule file where the teacher's schedule

is stored.

3.4.34 SUBROUTINE READ TEACHER SCHEDULE: A subroutine which reads a teacher's schedule from the teacher schedule file.

3.4.35 SUBROUTINE READ CLASS LIST: This subroutine reads the class list of a given class. The student numbers of all students enrolled in the class are stored as the vector LST.

3.4.36 SUBROUTINE INPUT AND WRITE GRADES: This subroutine inputs student grades and comments and writes the grades and comments to the student schedule file.

3.4.37 SUBROUTINE UPDATE STUDENT SCHEDULE: This subroutine updates the student schedule stored on the student schedule file and corrects the corresponding class size and class lists in the class information file.

3.4.38 SUBROUTINE MISSING PERSON: This subroutine displays a message that the required teacher, or student, is not currently defined in the data base.

- 3.4.39 SUBROUTINE UPDATE TEACHER SCHEDULE: This subroutine updates the teacher's schedule in the teacher schedule file and adds the teacher number to appropriate record of the class information file.
- 3.4.40 SUBROUTINE UPDATE GRADES: A subroutine to find the record number of the student schedule file for a given student, display the current grades, input corrected grades and comments, and store the corrected grades and comments in the student schedule file.
- 3.4.41 SUBROUTINE SET BYTE POINTERS: B1, B2, B3, and B4 are set to ten, thirteen, fifteen, and eighteen respectively to allow access to the mid-semester and end of semester grades and comments.
- 3.4.42 SUBROUTINE PRINT GRADES (ONE STUDENT): A subroutine to control inputting and processing of the grades for an individual student. The subroutine calls subroutines find student record, read grades, and returns control to the program menu.
- 3.4.43 SUBROUTINE READ GRADES: A subroutine to read the student schedule and existing grades from

the student schedule file.

- 3.4.44 SUBROUTINE PRINT MID-SEMESTER GRADES: A subroutine to print the schedule, teacher numbers, room numbers, and mid-semester grades.
- 3.4.45 SUBROUTINE PRINT FINIAL GRADES: A subroutine to print the student schedule, teacher number, room number, and end of semester grades and comments.
- 3.4.46 SUBROUTINE INPUT CLASS NUMBER: A subroutine which inputs a class number and searches the class index to find the record number of the class information file corresponding to the given class.
- 3.4.47 SUBROUTINE READ CLASS DATA: A subroutine to read the class number, teacher number, period number, class size, and class list from the class information file.
- 3.4.48 SUBROUTINE PRINT CLASS LIST: A subroutine to display the teacher number, period number, room number, class size, and the student numbers of all students currently enrolled in a given class.
- 3.4.49 SUBROUTINE PROCESS LISTS BY TEACHER: A subroutine which inputs a teacher number, finds the

corresponding record number of the teacher schedule file, and reads the teacher schedule from the file. This subroutine then calls the find class record, read class record, read class data, and print class list subroutines for each class.

3.4.50 SUBROUTINE PROCESS ALL CLASS LISTS: A subroutine to print all class lists. The subroutine calls subroutine in-order to produce a list of all teachers in order of increasing teacher number. The subroutine then reads each teacher's schedule and calls subroutines find class record, read class record, read class data, and print class lists for all classes defined in the data base.

3.4.51 SUBROUTINE CHANGE CLASS INDEX: A subroutine to copy the CINDX array into the INDX array to enable the balance routine to balance the class index structure.

3.4.52 SUBROUTINE PRINT GRADES (ALL STUDENTS): A subroutine to control the processing of all student grades. The subroutine first calls subroutine in-order to produce a list of all students in order of increasing student

number. Subroutines read student grades and print student grades are then called for all students contained in the data base.

3.4.53 SUBROUTINE UPDATE ROOM NUMBERS: A subroutine that inputs a class number, finds the existing room number, inputs a new room number and updates the class information file.

Program
SYSTEM MENU

programs

-Initialize files

-Build Schedules

-Retrieve Schedules

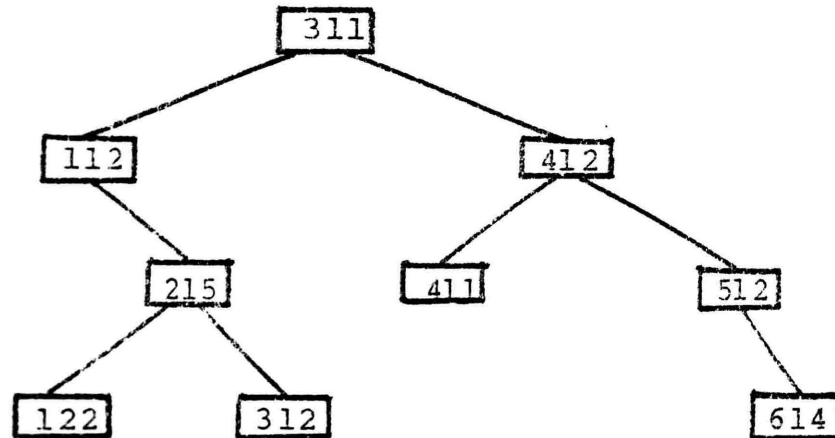
-Retrieve Class Lists

-Enter Grades

-Print Grades

-Update

Figure 14: Programs of the Student Information System



15-a Class-Index as a binary tree

Left pointer	Class Number	Right Pointer	Record Number
4	311	2	0
8	412	3	1
0	512	6	2
0	112	5	3
7	215	9	4
0	614	0	5
0	122	0	6
9	411	0	7
0	312	0	8

15-b Class-Index as an Array

Figure 15: Index Structure as a Binary Tree and an Array

```

10  REM PROCEDURE TO SEARCH BINARY TREE INDEX
20  PT = SP
30  IF PT = 0 THEN GOTO 110
40  IF NUMB < INDEX(PT,2) THEN IK = 1
50  IF NUMB > INDEX(PT,2) THEN IK = 3
60  IF NUMB = INDEX(PT,2) THEN GOTO 90
70  PT = INDEX(PT,IK)
80  GOTO 30
90  REC = INDEX(PT,4)
100 GOTO 120
110 PRINT "CLASS IS NOT IN THE INDEX"
120 RETURN

```

Line 20 Initializes the pointer PT to SP, the root of node of the index tree.

Line 30 Tests to determine if the current node is a terminal node of the tree. If the current node is a terminal node NUMB is not contained in the index.

Line 40 If NUMB is less than the current node set IK, a temporary pointer, to 1.

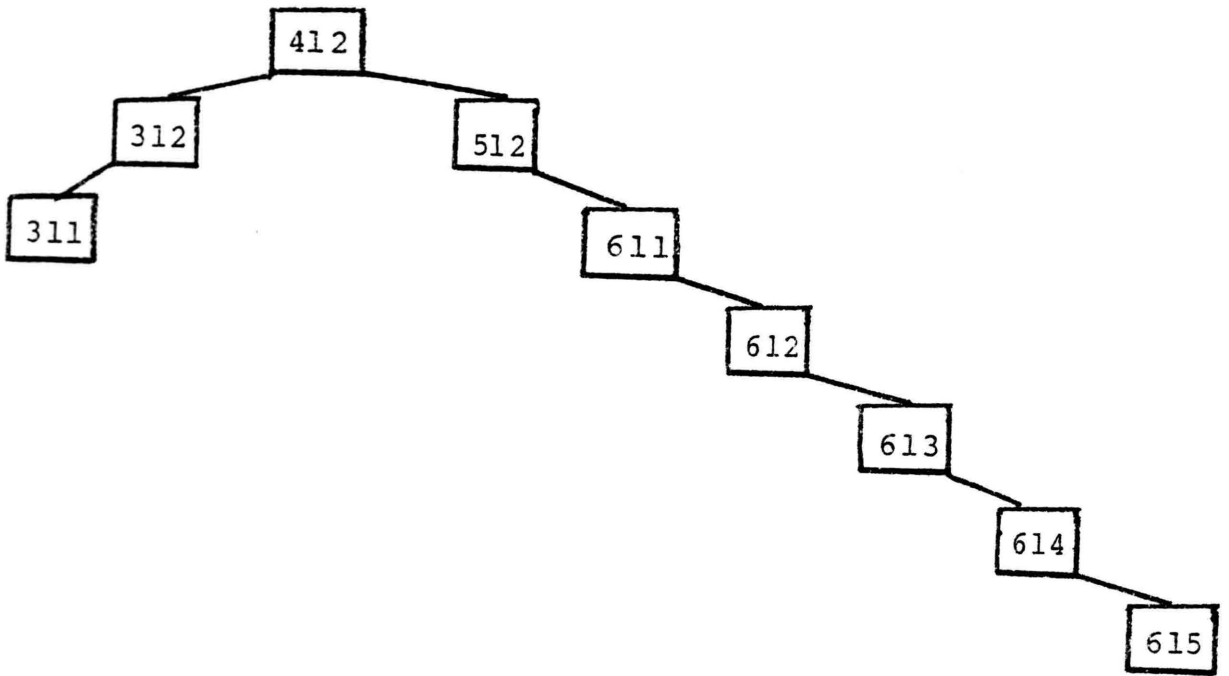
Line 50 If NUMB is greater than the current node set the temporary pointer to 3.

Line 60 If NUMB is equal to the current node the record has been found.

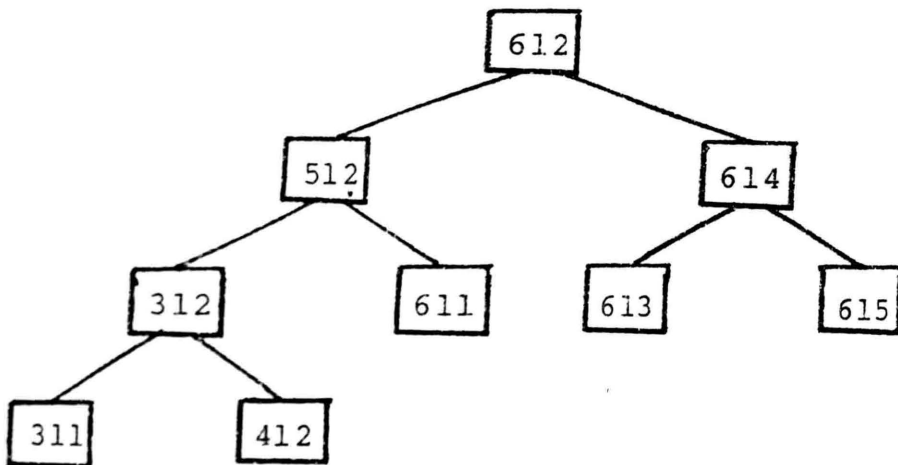
Line 70 Move down the tree by setting the pointer PT equal to the left or right pointer as determined above.

Line 80 Repeat until NUMB is located or a terminal node is located.

Figure 16: Basic code for searching an index structure



17-a Unbalanced Binary Tree Structure



17-b Balanced Binary Tree Structure

Figure 17: Binary Tree Structures

Step 1: Traverse in order and create the ODR vector:
 ODR(3 , 2 , 1 , 4 , 5 , 6 , 7 , 8 , 9)

Step 2: Calculate the maximum length of the full tree.
 $LF = 7 = (2^3 - 1)$

Step 3: Determine the number of elements of delete.
 $ND = 2 = (\text{length of index} - \text{length of full tree})$
 Create the OMIT vector and pack ODR.
 OMIT(1) = ODR(1) = 3
 OMIT(2) = ODR(3) = 1
 Pack ODR.
 ODR(2 , 4 , 5 , 6 , 7 , 8 , 9)

Step 4a: Set all odd pointers to zero.
 $INDEX(ODR(I),1) = INDEX(ODR(I),3) = 0; I \text{ odd.}$

Step 4b: Position pointers in following passes:
 pass 1: $INDEX(ODR(2),1) = ODR(1) \dots INDEX(4,1) = 2$
 $INDEX(ODR(2),3) = ODR(3) \dots INDEX(4,3) = 5$
 $INDEX(ODR(6),1) = ODR(5) \dots INDEX(8,1) = 7$
 $INDEX(ODR(6),3) = ODR(7) \dots INDEX(8,3) = 9$
 pass 2: $INDEX(ODR(4),1) = ODR(2) \dots INDEX(6,1) = 4$
 $INDEX(ODR(4),3) = ODR(6) \dots INDEX(6,3) = 8$

Step 5: Place omitted nodes into the tree.
 $INDEX(ODR(1),1) = OMIT(1) \dots INDEX(2,1) = 3$
 $INDEX(ODR(1),3) = OMIT(2) \dots INDEX(2,3) = 1$

Step 6: Set pointers of the inserted nodes to zero.
 $INDEX(OMIT(1),1) = 0 \dots \dots \dots INDEX(3,1) = 0$
 $INDEX(OMIT(1),3) = 0 \dots \dots \dots INDEX(3,3) = 0$
 $INDEX(OMIT(2),2) = 0 \dots \dots \dots INDEX(1,1) = 0$
 $INDEX(OMIT(2),3) = 0 \dots \dots \dots INDEX(1,3) = 0$

Step 7: Define the root node:
 $\text{Root node} = ODR((7 + 1) / 2) = ODR(4) = 6$

Unbalanced index.				Balanced index.			
2	412	4	0	0	412	0	0
0	312	3	1	2	312	1	1
0	311	0	2	0	311	0	2
0	512	5	3	2	512	5	3
0	611	6	4	0	611	0	4
0	612	7	5	4	612	8	5
0	613	8	6	0	613	0	6
0	614	9	7	7	614	9	7
0	615	0	8	0	615	0	8

Figure 18: Balance Routine

Variable	Description
INDX(500,4)	Array to internally store the student index structure.
STACK(256,2)	Array used as a stack by the in-order subroutine.
OMIT(256)	Vector to temporarily store nodes deleted in the balance procedure.
ODR(256)	Vector to store the pointers to the nodes in increasing order.
SCHD(7)	Vector to internally store the student or teacher schedule.
CINDX(150,4)	Array to internally store the class index structure.
SNDX(500,4)	Array to internally store the student index structure.
TINDX(50,4)	Array to internally store the teacher index structure.
CREC(7)	Vector to store the record numbers of the class information file.
TCH(7)	Vector to store the teacher schedule.
GMT(7)	Vector to store the mid-semester grades.
GF(7)	Vector to store the end of semester grades.
CMT(7)	Vector to store the mid-semester teacher comments.
CF(7)	Vector to store the end of semester comments.
INFO(5)	Vector to store class information: course number, teacher, period number room number, and class size.
LST(35)	Vector to store class lists.
GDE(35)	Vector to store grades for a class.
COM(35)	Vector to store comments for a class.

Figure 19: Variable Documentation

Program	Subroutines
BUILD SCHEDULES:	
	-Read Student Index (3.4.2)
	-Read Teacher Index (3.4.1)
	-Read Class Index (3.4.3)
	-Change Schedule (3.4.4)
	-Display Schedule (3.4.5)
	-Change Number (3.4.6)
	-Display Schedule (3.4.5)
	-Change Classes (3.4.7)
	-Display Schedule (3.4.5)
	-Find Student Record (3.4.8)
	-Write Student Schedule (3.4.9)
	-Update Class Index (3.4.10)
	-Revise Class Info(student) (3.4.11)
	-Write Teacher Schedule (3.4.30)
	-Update Class Index (3.4.10)
	-Revise Class Info(teacher) (3.4.12)
	-Balance Index (3.4.14)
	-In-Order (3.4.15)
	-Push Stack (3.4.28)
	-Pop Stack (3.4.29)
	-Length (3.4.17)
	-Position Pointers (3.4.16)
	-Write Class Index (3.4.19)
	-Write Teacher Index (3.4.18)
	-Write Student Index (3.4.20)
	-Change Class Index (3.4.19)

Figure 20: Program Build Schedules (Structure)

Program:	Subroutines
Retrieve Schedules:	
	-Read Student Index (3.4.2)
	-Read Teacher Index (3.4.1)
	-Read Class Index (3.4.3)
	-Retrieve Student Schedule
	(one student) (3.4.21)
	-Find Class Record (3.4.22)
	-Read Teacher/Room number
	(3.4.23)
	-Print Student Schedule
	(3.4.24)
	-Retrieve Teacher Schedule
	(one teacher) (3.4.25)
	-Find Class Record (3.4.22)
	-Read Teacher/Room number
	(3.4.23)
	-Print Teacher Schedule
	(3.4.30)
	-Retrieve Student Schedules
	(3.4.26)
	-In Order (3.4.15)
	-Push Stack (3.4.28)
	-Pop Stack (3.4.29)
	-Find Class Record (3.4.22)
	-Read Teacher/Room number
	(3.4.23)
	-Print Student Schedule
	(3.4.24)
	-Retrieve Teacher Schedules.
	(2.4.27)
	-In Order (3.8.15)
	-Push Stack (3.4.28)
	-Pop Stack (3.4.29)
	-Find Class Record (3.4.22)
	-Print Teacher Schedule
	(3.4.30)

Figure 21: Program Retrieve Schedules (Structure)

Program:	Subroutines
Update	-Read Student Index (3.4.2)
	-Read Class Index (3.4.3)
	-Read Teacher Index (3.4.1)
	-Update Student Schedule (3.4.37)
	-Missing Person (3.4.38)
	-Find Class Record (3.4.22)
	-Update Teacher Schedule (3.4.39)
	-Missing Person (3.4.38)
	-Find Class Record (3.4.22)
	-Update Grades (3.4.40)
	-Update Room numbers (3.4.53)

Figure 22: Program Update (Structure)

Program	Subroutines
Enter Grades:	
	-Read Teacher Index (3.4.1)
	-Read Class Index (3.4.3)
	-Read Student Index (3.4.2)
	-Byte Pointers mid-semester (3.4.31)
	-Byte Pointers finial (3.4.32)
	-Find Teacher Record (3.4.33)
	-Read Teacher Schedule (3.4.34)
	-Find Class Record (3.4.42)
	-Find Student Records (3.4.8)
	-Input and write grades (3.4.36)

Figure 23: Program Enter Grades (Structure)

Program

Subroutines

Print Grades:

- Read Student Index (3.4.2)
- Read Class Index (3.4.3)
- Set Byte Pointers (3.4.41)
- Print Grades (one student)
 - (3.4.42)
 - Find Student record (3.4.8)
 - Print mid-semester grades
 - (3.4.44)
 - Print final grades (3.4.45)
- Print grades (all students)
 - (3.4.52)
 - In Order (3.4.15)
 - Push Stack (3.4.28)
 - Pop Stack (3.4.29)
 - Read Grades (3.4.43)
 - Print mid-semester grades
 - (3.4.44)
 - Print final grades (3.4.45)

Figure 24: Program Retrieve Grades (Structure)

Program	Subroutines
Retrieve Class Lists:	
	-Read Teacher Index (3.4.1)
	-Read Class Index (3.4.3)
	-Input Class Number (3.4.46)
	-Find Class Record (3.4.22)
	-Read Class Data (3.4.47)
	-Print Class Lists (3.4.48)
	-Process All Class Lists (3.4.50)
	-In Order (3.4.15)
	-Push Stack (3.4.28)
	-Pop Stack (3.4.29)
	-Find Class Record (3.4.22)
	-Read Class Data (3.4.47)
	-Print Class List (3.4.48)

Figure 25: Program Retrieve Class Lists (Structure)

CHAPTER 1V

CONCLUSION AND REMARKS

The purpose of this study was to investigate the feasibility of using microcomputers to maintain a student information system to serve a small high school. The memory limitations, both core and direct access, and the time required to retrieve information where the primary concerns in the system's design and implementation. To discuss the results of the study it is first necessary to discuss the specifications of the specific micro computer used in the implementation of the system developed. The system developed in this study was implemented on an Apple II PLUS micro computer with 48K random access memory and external storage in the form of five and a quarter inch magnetic floppy diskettes. The 48K RAM proved to be sufficient to execute the system software and the maximum number of students the system can maintain is dependent upon the external storage limitations. To analysis the maximum capacity of the system we must first analysis the storage capacity of a floppy diskette when initialized on an Apple II micro computer.

When the Apple's disk operating system initializes a

diskette, the diskette is divided into thirty five tracks with each track containing sixteen sectors. Each sector will store two hundred and fifty six bytes of data and four tracks are reserved to store the disk operating system and disk directory. This leaves thirty one tracks or four hundred and ninety six sectors available to the user to store software or data files. The number of sectors required to store a file is given in the catalog listing of the directory of the diskette. The catalog listing of the diskette which stores the system software for the system developed in this study is given in figure 26 and shows that a total of one hundred and seventy three sectors are required to store the system software. This leaves three hundred and twenty three sectors available to store the student schedule, teacher schedule, class information, and the three index files. Figure 27(a) gives the disk directory when the six data files contain the data for three hundred students, fifty teachers, and one hundred and twenty classes. As figure 27(a) illustrates, the system developed in this study will maintain a student information system for three hundred students when implemented on an Apple II micro computer with 48K RAM and a single disk drive. Further analysis of the unused forty eight sectors on the disk imply that the student information system designed in this study

can accommodate a maximum of three hundred and fifty students when implemented on a one disk drive Apple II PLUS micro computer configuration.

The addition of a second disk drive greatly increases the capacity of the system. If the system software is stored on diskette number one, operating in drive one, and the data files are stored on diskette number two, operating in drive two, the maximum number of students increases to over five hundred. Figure 27(b) shows the catalog listing when the data files contain the data for five hundred students, fifty teachers, and one hundred and fifty courses. A total of three hundred and ninety six sectors are required to store the data files which leaves one hundred unused sectors. Further analysis of the available space indicates that a two disk drive Apple II PLUS with 48K RAM hardware configuration, will maintain a data base for six hundred and fifty students. An added advantage of using a dual disk drive system is that the three hundred and twenty three unused sectors on diskette number one could be utilized to store the student and teacher name relations.

Three hundred and twenty three sectors represents approximately 82K bytes of storage. If each name is restricted to thirty characters, each record (number,name) would require thirty seven bytes when end of field and

record markers are included. Since the analysis above gives the maximum size of the system as six hundred and fifty students with sixty five teachers, the name relations would require a total of approximately 30K bytes. Thus diskette number one would accommodate the files required to store the name relations. The addition of these relations would allow student schedules, teacher schedules, and student grade reports to be printed by name.

Although the above analysis of the direct access storage capabilities of the Apple II indicate that a maximum of six hundred and fifty students could be maintained with this system, problems do occur with the software. The system, as designed, will accommodate a maximum of five hundred students since the program update files, listed in appendix V, utilizes the complete 48K of random access memory. The system, as presented in this study, would have to be slightly modified to accommodate more than five hundred students.

In conclusion, this investigation did find that the microcomputer is a viable alternative for maintaining a small stand-alone data base. It is indeed feasible for a school, enrollment around five hundred students, to consider implementing its student information system on a dual disk drive micro computer hardware configuration. The student

information system presented in this study will maintain the data base required for five hundred students and slight modifications of the system design would increase the maximum number of students to six hundred and fifty. This study further indicates that student and teacher name relations could be added to the system to allow for schedule and grade reports to be printed by teacher and student name. Further modifications should be made to the system software before it is implemented. The system contains all the information needed to assist in the scheduling process of a school. The system could easily be modified to allow student and teacher classes to be entered without regard to period. The system could supply the period and room numbers and report schedule conflicts.

A	036	BUILD SCHEDULES
A	020	RETRIEVE CLASS LISTS
A	005	MENU
A	026	RETRIEVE SCHEDULES
A	020	ENTER GRADES
A	024	PRINT GRADES
A	036	UPDATE FILES
A	006	INITIALIZE DATA FILES

Total of 173 sectors required to store the system software.

Figure 26: Storage requirement for the system software

T 025 STUDENT-INDEX
T 005 TEACHER-INDEX
T 008 CLASS-INDEX
T 008 TEACHER-SCHEDULE
T 106 CLASS-INFORMATION
T 123 STUDENT-SCHEDULE

27-a Files initialized for 300 students (275 sectors)

T 033 STUDENT-INDEX
T 005 TEACHER-INDEX
T 010 CLASS-INDEX
T 008 TEACHER-SCHEDULES
T 134 CLASS-INFORMATION
T 206 STUDENT-SCHEDULE

27-b Files initialized for 500 students (396 sectors)

Figure 27: Data files storage requirements

APPENDIX A

PROGRAM SYSTEM MENU

```
100 HOME : PRINT : PRINT
110 D$ = CHR$ (4)
120 PRINT TAB( 7) "STUDENT INFORMATION SYSTEM"
130 PRINT
140 PRINT TAB( 15) "CREATED BY"
150 PRINT
160 PRINT TAB( 13) "DOUG WAECHTER"
170 PRINT : PRINT : PRINT
180 PRINT " MENU"
190 PRINT : PRINT
200 PRINT " 1.....BUILD SCHEDULES"
210 PRINT " 2.....RETRIEVE SCHEDULES"
220 PRINT " 3.....RETRIEVE CLASS LISTS"
230 PRINT " 4.....ENTER GRADES"
240 PRINT " 5.....PRINT GRADES"
250 PRINT " 6.....UPDATE FILES"
260 PRINT " 7.....EXIT SYSTEM"
270 PRINT : PRINT
280 INPUT "ENTER THE NUMBER OF YOUR CHOICE ";ANS
290 IF ANS < 1 OR ANS > 7 THEN GOTO 100
300 HOME : PRINT : PRINT
310 PRINT TAB( 13) "LOADING PROGRAM"
320 PRINT : PRINT
330 PRINT TAB( 14) "PLEASE WAIT"
340 IF ANS = 1 THEN PRINT D$; "RUN BUILD SCHEDULES"
350 IF ANS = 2 THEN PRINT D$; "RUN RETRIEVE SCHEDULES"
360 IF ANS = 3 THEN PRINT D$; "RUN RETRIEVE CLASS LISTS"
370 IF ANS = 4 THEN PRINT D$; "RUN ENTER GRADES"
380 IF ANS = 5 THEN PRINT D$; "RUN PRINT GRADES"
390 IF ANS = 6 THEN PRINT D$; "RUN UPDATE FILES"
400 IF ANS = 7 THEN GOTO 420
410 GOTO 100
420 HOME
430 END
```

APPENDIX B

PROGRAM INITIALIZE FILES

```
100 REM PROGRAM TO INITIALIZE FILES.
110 REM
120 D$ = CHR$ (4)
130 PRINT D$; "OPEN STUDENT-INDEX"
140 PRINT D$; "WRITE STUDENT-INDEX"
150 PRINT 0: PRINT 0: PRINT 0
160 PRINT D$; "OPEN TEACHER-INDEX"
170 PRINT D$; "WRITE TEACHER-INDEX"
180 PRINT 0: PRINT 0: PRINT 0
190 PRINT D$; "OPEN CLASS-INDEX"
200 PRINT D$; "WRITE CLASS-INDEX"
210 PRINT 0: PRINT 0: PRINT 0
220 PRINT D$; "CLOSE"
230 PRINT D$; "OPEN TEACHER-SCHEDULE,L31"
240 FOR I = 1 TO 65
250 I1 = I - 1
260 PRINT D$; "WRITE TEACHER-SCHEDULE,R";I1; ",B0"
270 PRINT 0
280 FOR K = 3 TO 27 STEP 4
290 PRINT D$; "WRITE TEACHER-SCHEDULE,R";I1; ",B";K
300 PRINT 0
310 NEXT K
320 NEXT I
330 PRINT D$; "CLOSE"
340 REM
350 PRINT D$; "OPEN CLASS-INFO,L225"
360 FOR I = 1 TO 150
370 I1 = I - 1
380 PRINT D$; "WRITE CLASS-INFO,R";I1; ",B0"
390 PRINT 0
400 PRINT D$; "WRITE CLASS-INFO,R";I1; ",B4"
410 PRINT 0
420 PRINT D$; "WRITE CLASS-INFO,R";I1; ",B7"
430 PRINT 0
440 PRINT D$; "WRITE CLASS-INFO,R";I1; ",B9"
450 PRINT 0
460 PRINT D$; "WRITE CLASS-INFO,R";I1; ",B12"
470 PRINT 0
```



```
480 NEXT I
490 PRINT D$;"CLOSE"
500 D$ = CHR$ (4)
510 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
520 FOR I = 1 TO 500
530 BYT = 0
540 GOSUB 700
550 FOR K = 1 TO 7
560 BYT = 6 + 14 * (K - 1)
570 GOSUB 700
580 BYT = 10 + 14 * (K - 1)
590 GOSUB 700
600 BYT = 13 + 14 * (K - 1)
610 GOSUB 700
620 BYT = 15 + 14 * (K - 1)
630 GOSUB 700
640 BYT = 18 + 14 * (K - 1)
650 GOSUB 700
660 NEXT K
670 NEXT I
680 PRINT D$;"CLOSE STUDENT-SCHEDULE"
690 END
700 REM SUBROUTINE
710 PRINT D$;"WRITE STUDENT-SCHEDULE,R";I - 1;","B";BYT
720 PRINT 0
730 RETURN
```

APPENDIX C

PROGRAM BUILD SCHEDULES

```

100 REM PROGRAM BUILD-SCHEDULES
110 REM
120 DIM INDXS(500,4)
130 DIM STACK(400,2)
140 DIM OMIT(256)
150 DIM ODR(500)
160 DIM SCHD(7)
170 DIM CINDX(150,4)
180 DS = CHR$(4)
190 HOME : PRINT : PRINT
200 PRINT " THIS PROGRAM BUILDS BOTH"
210 PRINT
220 PRINT "THE STUDENT-SCHEDULE AND THE "
230 PRINT : PRINT " THE TEACHER-SCHEDULE FILES."
240 PRINT : PRINT : PRINT "ENTER THE NUMBER OF YOUR CHOICE"
250 PRINT : PRINT : PRINT
260 PRINT "1..FOR STUDENT-SCHEDULE.."
270 PRINT : PRINT
280 PRINT "2..FOR TEACHER-SCHEDULE.."
290 PRINT : PRINT : INPUT ANS
300 GOSUB 1470
310 IF ANS = 1 THEN GOSUB 1190
320 IF ANS = 2 THEN GOSUB 1340
330 IF ANS < > 1 AND ANS < > 2 THEN GOTO 190
340 IF ANS = 1 THEN N$ = "STUDENT"
350 IF ANS = 2 THEN N$ = "TEACHER"
360 GOSUB 1470
370 HOME : PRINT : PRINT
380 PRINT "THIS PROGRAM BUILDS THE ";N$;"-SCHEDULE";
385 PRINT " AND THE ";N$;"-INDEX FILES."
390 PRINT : PRINT : PRINT
400 PRINT "YOU CAN TERMINATE THIS PROGRAM AT ANY"
410 PRINT
420 PRINT "TIME AND THE DATA WILL BE SAVED"
430 PRINT : PRINT : PRINT
440 PRINT "SIMPLY PRESS THE 'E' KEY AT THE PROMPT"
450 PRINT : PRINT : PRINT
460 PRINT " <PRESS 'E' TO EXIT> "
```

```

470 PRINT : PRINT : PRINT
480 PRINT " PRESS THE RETURN KEY TO CONTINUE ":INPUT Z$
490 HOME : PRINT : PRINT : PRINT : PRINT : PRINT
500 PRINT "ENTER 'E' TO EXIT.....": PRINT
510 INPUT "RETURN TO CONTINUE.....";A$
520 IF A$ = "E" THEN GOTO 720
530 HOME
540 PRINT : PRINT : PRINT "ENTER THE ";N$;" NUMBER ";
550 INPUT STUNO
560 PRINT : PRINT "ENTER COURSES BY PERIOD : "
570 PRINT
580 FOR I = 1 TO 7
590 PRINT : PRINT "PERIOD ";I;" ";
600 INPUT SCHD(I)
610 NEXT I
620 PRINT
630 PRINT "DO YOU WANT TO MAKE ANY CHANGES"
640 INPUT "....ENTER (Y/N) ...";Z$
650 IF Z$ = "Y" THEN GOSUB 2500
660 NST = NST + 1
670 GOSUB 900
680 IF ANS = 1 THEN GOSUB 1040
690 IF ANS = 2 THEN GOSUB 3510
700 GOTO 490
710 REM
720 GOSUB 760
730 PRINT D$;"RUN MENU,D1"
740 REM *****
750 REM SUBROUTINE EXIT
760 REM *****
770 PT = SS
780 IF ANS = 1 AND NST - OBAL > 25 THEN GOSUB 3750
790 IF ANS = 1 THEN GOSUB 3960
800 IF ANS = 2 AND NST - OBAL > 10 THEN GOSUB 3750
810 IF ANS = 2 THEN GOSUB 4080
820 GOSUB 3660
830 FLAG = 1
840 IF NCLASS - OCBAL > 5 THEN GOSUB 3750
850 GOSUB 3850
860 RETURN
870 REM *****
880 REM SUBROUTINE FIND STUDENT-RECORD
890 REM *****
900 PT = SS
910 IF PT < > 0 THEN GOTO 940
920 PT = 1:SS = 1

```

```

930 GOTO 1010
940 IF STNO < INDXS(PT,2) THEN IK = 1
950 IF STNO > INDXS(PT,2) THEN IK = 3
960 IF STNO = INDXS(PT,2) THEN RETURN
970 LAST = PT
980 PT = INDXS(PT,IK)
990 IF PT < > 0 THEN GOTO 940
1000 INDXS(LAST,IK) = NST
1010 INDXS(NST,2) = STNO
1020 INDXS(NST,4) = NST - 1
1030 RETURN
1040 REM *****
1050 REM SUBROUTINE WRITE STUDENT-SCHEDULE
1060 REM *****
1070 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
1080 PRINT D$;"WRITE STUDENT-SCHEDULE,R";NST - 1
1090 PRINT STNO
1100 FOR I = 1 TO 7
1110 BYT = 6 + (I - 1) * 14
1120 PRINT D$;"WRITE STUDENT-SCHEDULE,R";NST - 1;"B";BYT
1130 PRINT SCHD(I)
1140 NEXT I
1150 PRINT D$;"CLOSE STUDENT-SCHEDULE"
1160 GOSUB 3000
1170 RETURN
1180 REM *****
1190 REM SUBROUTINE READ STUDENT-INDEX
1200 REM *****
1210 REM
1220 PRINT D$;"OPEN STUDENT-INDEX"
1230 PRINT D$;"READ STUDENT-INDEX"
1240 INPUT NST: INPUT SS: INPUT OBAL
1250 IF NST = 0 THEN GOTO 1310
1260 FOR I = 1 TO NST
1270 FOR J = 1 TO 4
1280 INPUT INDXS(I,J)
1290 NEXT J
1300 NEXT I
1310 PRINT D$;"CLOSE STUDENT-INDEX"
1320 RETURN
1330 REM *****
1340 REM SUBROUTINE READ TEACHER-INDEX
1350 REM *****
1360 PRINT D$;"OPEN TEACHER-INDEX"
1370 PRINT D$;"READ TEACHER-INDEX"
1380 INPUT NST: INPUT SS: INPUT OBAL
1390 IF NST = 0 THEN GOTO 1440

```

```

1400 FOR I = 1 TO NST
1410 FOR J = 1 TO 4
1420 INPUT INDX(I,J)
1430 NEXT J: NEXT I
1440 PRINT D$; "CLOSE TEACHER-INDEX"
1450 RETURN
1460 REM *****
1470 REM SUBROUTINE READ CLASS-INDEX
1480 REM *****
1490 PRINT D$; "OPEN CLASS-INDEX,D2"
1500 PRINT D$; "READ CLASS-INDEX"
1510 INPUT NCLASS: INPUT S2C: INPUT OCBAL
1520 IF NCLASS = 0 THEN GOTO 1580
1530 FOR I = 1 TO NCLASS
1540 FOR J = 1 TO 4
1550 INPUT CINDX(I,J)
1560 NEXT J
1570 NEXT I
1580 PRINT D$; "CLOSE CLASS-INDEX"
1590 RETURN
1600 REM *****
1610 REM SUBROUTINE IN-ORDER
1620 REM *****
1630 PT = SS
1640 IF INDXS(PT,1) < > 0 THEN GOTO 1760
1650 LN = LN + 1
1660 ODR(LN) = PT
1670 D = 1
1680 GOSUB 1850
1690 D = 0
1700 PT = INDXS(PT,3)
1710 IF PT < > 0 THEN GOTO 1640
1720 GOSUB 1920
1730 IF D = 3 THEN GOTO 1810
1740 IF D = 1 THEN GOTO 1720
1750 GOTO 1650
1760 D = 2
1770 GOSUB 1850
1780 D = 0
1790 PT = INDXS(PT,1)
1800 GOTO 1640
1810 RETURN
1820 REM *****
1830 REM SUBROUTINE PUSH
1840 REM *****
1850 LST = LST + 1
1860 STACK(LST,1) = PT

```

```

1870  STACK(LST,2) = D
1880  RETURN
1890  REM *****
1900  SUBROUTINE POP
1910  REM *****
1920  IF LST = 0 THEN GOTO 1970
1930  PT = STACK(LST,1)
1940  D = STACK(LST,2)
1950  LST = LST - 1
1960  RETURN
1970  D = 3
1980  RETURN
1990  REM *****
2000  REM SUBROUTINE LENGTH
2010  REM *****
2020  K1 = 1
2030  IF 2 ^ K1 > LN THEN GOTO 2050
2040  K1 = K1 + 1: GOTO 2030
2050  EX = K1 - 1
2060  LF = 2 ^ EX - 1
2070  ND = LN - LF
2080  K2 = 0
2090  FOR K1 = 1 TO 500 STEP 2
2100  K2 = K2 + 1
2110  OMIT(K2) = ODR(K1)
2120  IF K2 = ND GOTO 2140
2130  NEXT K1
2140  K5 = 2 * ND
2150  FOR K3 = 1 TO LN
2160  IF K3 > ND THEN GOTO 2190
2170  ODR(K3) = ODR(2 * K3)
2180  GOTO 2210
2190  K5 = K5 + 1
2200  ODR(K3) = ODR(K5)
2210  NEXT K3
2220  RETURN
2230  REM *****
2240  REM SUBROUTINE ENTER-POINTERS
2250  REM *****
2260  FOR K1 = 1 TO LF STEP 2
2270  INDX(ODR(K1),1) = 0
2280  INDX(ODR(K1),3) = 0
2290  NEXT K1
2300  PASS = 1
2310  KS = PASS + 1
2320  KB = 2 ^ PASS
2330  FOR I = KB TO LF STEP 2 ^ KS

```

```

2340  INDX(ODR(I),1) = ODR(I - 2 ^ (PASS - 1))
2350  INDX(ODR(I),3) = ODR(I + 2 ^ (PASS - 1))
2360  NEXT I
2370  PASS = PASS + 1
2380  IF 2 ^ PASS < LF THEN GOTO 2310
2390  KD = 1
2400  INDXS(ODR(KD),1) = OMIT(KD)
2410  INDXS(ODR(KD),3) = OMIT(KD + 1)
2420  KD = KD + 2
2430  IF KD < = ND THEN GOTO 2400
2440  SS = ODR(2 ^ (EX - 1))
2450  FOR KZ = 1 TO ND
2460  INDXS(OMIT(KZ),1) = 0
2470  INDXS(OMIT(KZ),3) = 0
2480  NEXT KZ
2490  RETURN
2500  REM *****
2510  REM SUBROUTINE CHANGE-SCHEDULE
2520  REM *****
2530  PRINT : PRINT
2540  GOSUB 2660
2550  PRINT "DO YOU WANT TO CHANGE THE"
2560  PRINT "      ";N$;" NUMBER (Y/N) "
2570  INPUT H$
2580  IF H$ = "Y" THEN GOSUB 2890
2590  IF H$ = "N" THEN GOSUB 2800
2600  IF H$ < > "Y" AND H$ < > "N" THEN GOTO 2550
2610  PRINT : PRINT "DO YOU WANT TO MAKE ANY MORE CHANGES"
2615  PRINT "      ";: INPUT "(Y/N) ";H$
2620  IF H$ = "Y" THEN GOTO 2530
2630  IF H$ = "N" THEN RETURN
2640  IF H$ < > "Y" AND H$ < > "N" THEN GOTO 2610
2650  GOTO 2610
2660  REM *****
2670  REM SUBROUTINE DISPLAY SCHEDULE
2680  REM *****
2690  HOME
2700  PRINT "SCHEDULE GIVEN IS: "
2710  PRINT
2720  PRINT "      ";N$;" NUMBER      ";STN
2730  PRINT
2740  PRINT "PERIOD      COURSE"
2750  FOR I = 1 TO 7
2760  PRINT
2770  PRINT I,SCHD(I)
2780  NEXT I: PRINT
2790  RETURN

```

```

2795 REM
2800 REM *****
2810 REM SUBROUTINE CHANGE CLASSES
2820 REM *****
2830 GOSUB 2660
2840 PRINT : PRINT
2845 PRINT "ENTER PERIOD NUMBER, COURSE NUMBER"
2850 INPUT J, COU
2860 SCHD(J) = COUD
2870 GOSUB 2660
2880 RETURN
2890 REM *****
2900 REM SUBROUTINE CHANGE-NUMBER
2910 REM *****
2920 HOME : PRINT : PRINT : PRINT : PRINT
2930 PRINT "OLD ", N$, " NUMBER IS: "; STUNO: PRINT
2940 PRINT : PRINT : PRINT : PRINT
2950 PRINT : PRINT : PRINT : PRINT
2960 PRINT "ENTER NEW ", N$, "NUMBER ";
2970 INPUT STNO
2980 GOSUB 2660
2990 RETURN
3010 REM *****
3020 REM SUBROUTINE UPDATE-CLASS-INFO
3030 REM *****
3040 PRINT D$; "OPEN CLASS-INFO, L225"
3050 FOR J = 1 TO 7
3060 PT = S2C
3070 IF PT < > 0 THEN 3110
3080 S2C = 1: CSIZE = 1
3090 PT = 1
3100 NCLASS = 1: GOTO 3190
3110 IF SCHD(J) < CINDX(PT, 2) THEN IK = 1
3120 IF SCHD(J) > CINDX(PT, 2) THEN IK = 3
3130 IF SCHD(J) = CINDX(PT, 2) THEN GOTO 3260
3140 LAST = PT
3150 PT = CINDX(PT, IK)
3160 IF PT < > 0 THEN 3110
3170 NCLASS = NCLASS + 1
3180 CINDX(LAST, IK) = NCLASS
3190 CINDX(NCLASS, 2) = SCHD(J)
3200 CINDX(NCLASS, 4) = NCLASS - 1
3210 PRINT D$; "WRITE CLASS-INFO, R"; NCLASS - 1
3220 PRINT SCHD(J)
3230 PRINT D$; "WRITE CLASS-INFO, R"; NCLASS - 1; ", B7"
3240 PRINT J
3250 PT = NCLASS

```



```

3260 REC = CINDX (PT,4)
3270 IF ANS = 1 THEN GOSUB 3320
3280 IF ANS = 2 THEN GOSUB 3450
3290 NEXT J
3300 PRINT D$; "CLOSE CLASS-INFO"
3310 RETURN
3320 REM *****
3330 REM SUBROUTINE REVISE CLASS-INFO
3340 REM (STUDENT)
3350 REM *****
3360 PRINT D$; "READ CLASS-INFO,R";REC;"",B12"
3370 INPUT CSIZE
3380 CSIZE = CSIZE + 1
3390 PRINT D$; "WRITE CLASS-INFO,R";REC;"",B12"
3400 PRINT CSIZE
3410 BYT = 15 + (CSIZE - 1) * 6
3420 PRINT D$; "WRITE CLASS-INFO,R";REC;"",B";BYT
3430 PRINT STNO
3440 RETURN
3450 REM *****
3460 REM SUBROUTINE REVISE CLASS-INFO (TEACHER)
3470 REM *****
3480 PRINT D$; "WRITE CLASS-INFO,R";REC;"",B4"
3490 PRINT STNO
3500 RETURN
3510 REM *****
3520 REM SUBROUTINE WRITE TEACHER-SCHEDULE
3530 REM *****
3540 PRINT D$; "OPEN TEACHER-SCHEDULE,L31"
3550 REM
3560 PRINT D$; "WRITE TEACHER-SCHEDULE,R";NST - 1
3570 PRINT STNO
3580 FOR I = 1 TO 7
3590 BYT = 3 + (I - 1) * 4
3600 PRINT D$; "WRITE TEACHER-SCHEDULE,R";NST - 1",B";BYT
3610 PRINT SCHD(I)
3620 NEXT I
3630 PRINT D$; "CLOSE TEACHER-SCHEDULE"
3640 GOSUB 3000
3650 RETURN
3660 REM *****
3670 REM SUBROUTINE CHANGE CLASS INDEX)
3680 REM *****
3690 FOR I = 1 TO NCLASS
3700 FOR J = 1 TO 4
3710 INDX(I,J) = CINDX(I,J)
3720 NEXT J: NEXT I

```

```

3730  SS = S2C
3740  RETURN
3750  REM *****
3760  REM  SUBROUTINE BALANCE-INDEX
3770  REM *****
3780  GOSUB 1610
3790  GOSUB 2000
3800  GOSUB 2240
3810  IF ANS = 1 OR ANS = 2 THEN OBAL = NST
3820  IF FLAG = 1 THEN OCBAL = NCLASS
3830  RETURN
3840  REM *****
3850  REM  SUBROUTINE WRITE CLASS-INDEX
3860  REM *****
3870  PRINT D$;"OPEN CLASS-INDEX"
3880  PRINT D$;"WRITE CLASS-INDEX"
3890  PRINT NCLASS: PRINT SS: PRINT OCBAL
3900  FOR I = 1 TO NCLASS
3910  FOR J = 1 TO 4
3920  PRINT INDX(I,J)
3930  NEXT J: NEXT I
3940  PRINT D$;"CLOSE CLASS-INDEX"
3950  RETURN
3960  REM *****
3970  REM  SUBROUTINE WRITE STUDENT-INDEX"
3980  REM *****
3990  PRINT D$;"OPEN STUDENT-INDEX"
4000  PRINT D$;"WRITE STUDENT-INDEX"
4010  PRINT NST: PRINT SS: PRINT OBAL
4020  FOR I = 1 TO NST
4030  FOR J = 1 TO 4
4040  PRINT INDX(I,J)
4050  NEXT J: NEXT I
4060  PRINT D$;"CLOSE STUDENT-INDEX"
4070  RETURN
4080  REM *****
4090  REM  SUBROUTINE WRITE TEACHER-INDEX
4100  REM *****
4110  PRINT D$;"OPEN TEACHER-INDEX"
4120  PRINT D$;"WRITE TEACHER-INDEX"
4130  PRINT NST: PRINT SS: PRINT OBAL
4140  FOR I = 1 TO NST
4150  FOR J = 1 TO 4
4160  PRINT INDX(I,J)
4170  NEXT J: NEXT I
4180  PRINT D$;"CLOSE TEACHER-INDEX"
4190  RETURN

```

APPENDIX D

PROGRAM RETRIEVE SCHEDULES LISTING

```
45  REM
50  REM  PROGRAM RETRIEVE SCHEDULES
52  REM
55  REM
60  DIM SNDX(500,4)
65  DIM TINDX(50,4)
70  DIM CINDX(150,4)
75  DIM CREC(7)
80  DIM RM(7)
85  DIM TCH(7)
90  DIM SCHC(7)
95  DIM ODR(500)
105 DIM STACK(256,2)
110 D$ = CHR$(4)
115 GOSUB 235
120 GOSUB 295
125 GOSUB 360
130 HOME : PRINT : PRINT
135 PRINT " THIS PROGRAM DISPLAYS EITHER "
140 PRINT : PRINT
145 PRINT "1.....STUDENT SCHEDULE..."
150 PRINT : PRINT
155 PRINT "2.....TEACHER SCHEDULE..."
160 PRINT : PRINT
165 PRINT "3.....TO PRINT ALL STUDENT SCHEDULES"
170 PRINT : PRINT
175 PRINT "4.....TO PRINT ALL TEACHER SCHEDULES"
180 PRINT : PRINT
185 PRINT "5.....TO END THE PROGRAM"
190 PRINT : PRINT
195 PRINT " ENTER THE NUMBER OF YOUR CHOICE...."
200 PRINT : PRINT : PRINT : INPUT ANS
205 IF ANS = 1 THEN GOSUB 425
210 IF ANS = 2 THEN GOSUB 560
215 IF ANS = 3 THEN GOSUB 1080
220 IF ANS = 4 THEN GOSUB 1620
225 IF ANS = 5 THEN GOTO 233
230 GOTO 130
```

```

233 REM TERMINATE EXECUTION AND RUN MENU PROGRAM
234 PRINT D$;"RUN MENU,D1"
235 REM
240 REM *****
242 REM SUBROUTINE READ STUDENT INDEX
243 REM*****
245 PRINT D$;"OPEN STUDENT-INDEX,D2"
250 PRINT D$;"READ STUDENT-INDEX"
255 INPUT NST: INPUT SP: INPUT OBAL
260 FOR I = 1 TO NST
265 FOR J = 1 TO 4
270 INPUT SNDX(I,J)
275 NEXT J
280 NEXT I
285 PRINT D$;"CLOSE STUDENT-INDEX"
290 RETURN
293 REM *****
295 REM SUBROUTINE READ TEACHER INDEX
300 REM *****
305 PRINT D$;"OPEN TEACHER-INDEX"
310 PRINT D$;"READ TEACHER-INDEX"
315 INPUT NTEA: INPUT TP: INPUT OBAL
320 IF NTEA = 0 THEN GOTO 350
325 FOR I = 1 TO NTEA
330 FOR J = 1 TO 4
335 INPUT TINDX(I,J)
340 NEXT J
345 NEXT I
350 PRINT D$;"CLOSE"
352 RETURN
353 REM *****
360 REM SUBROUTINE READ CLASS INDEX
365 REM *****
370 PRINT D$;"OPEN CLASS-INDEX"
375 PRINT D$;"READ CLASS-INDEX"
380 INPUT NCLASS: INPUT CP: INPUT OCBAL
385 IF NCLASS = 0 THEN GOTO 415
390 FOR I = 1 TO NCLASS
395 FOR J = 1 TO 4
400 INPUT CINDX(I,J)
405 NEXT J
410 NEXT I
415 PRINT D$;"CLOSE"
420 RETURN
422 REM *****
425 REM SUBROUTINE RETREIVE STUDENT

```

```

426 REM                                     (ONE STUDNENT)
430 REM *****
435 HOME : PRINT : PRINT
440 PRINT "  ENTER STUDENT NUMBER"
445 PRINT : PRINT : INPUT NUMB
448 SNUMB = NUMB
450 PT = SP
455 IF PT = 0 THEN GOTO 481
460 IF NUMB < SNDX(PT,2) THEN IK = 1
465 IF NUMB > SNDX(PT,2) THEN IK = 3
470 IF NUMB = SNDX(PT,2) THEN GOTO 488
475 PT = SNDX(PT,IK)
480 GOTO 455
481 PRINT : PRINT
482 PRINT "STUDENT      ";NUMB
483 PRINT : PRINT "IS NOT IN THE DATA BASE"
485 PRINT : PRINT
486 INPUT "PRESS RETURN TO CONTINUE  ";ZZ$
487 GOTO 550
488 REC = SNDX(PT,4)
490 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
495 FOR I = 1 TO 7
500 BYT = 6 + (I - 1) * 14
505 PRINT D$;"READ STUDENT-SCHEDULE,R";REC;"B";BYT
510 INPUT SCHD(I)
515 NEXT I
520 PRINT D$;"CLOSE"
525 FOR M = 1 TO 7
530 CNUMB = SCHD(M)
535 PT = CP
540 GOSUB 1430
542 CREC(M) = REC
544 NEXT M
546 GOSUB 1500
548 GOSUB 1560
550 RETURN
555 REM *****
560 REM SUBROUTINE FIND TEACHER RECORD
565 REM *****
570 HOME : PRINT : PRINT
575 PRINT "  ENTER TEACHER NUMBER  "
580 PRINT : PRINT : PRINT : INPUT NUMB
582 TNUMB = NUMB
585 PT = TP
590 IF PT = 0 THEN GOTO 615
595 IF NUMB < TINDX(PT,2) THEN IK = 1
600 IF NUMB > TINDX(PT,2) THEN IK = 3

```

```

605 IF NUMB = TINDX(PT,2) THEN GOTO 625
610 PT = TINDX(PT,IK)
612 GOTO 590
615 PRINT:PRINT"TEACHER ";NUMB;" IS NOT IN THE DATA BASE"
618 PRINT : PRINT : INPUT "PRESS RETURN TO CONTINUE ";ZZ$
620 GOTO 710
625 REC = TINDX(PT,4)
628 PRINT D$;"OPEN TEACHER-SCHEDULE,L31"
630 FOR I = 1 TO 7
635 BYT = 3 + (I - 1) * 4
640 PRINT D$;"READ TEACHER-SCHEDULE,R";REC;" ,B";BYT
645 INPUT SCHD(I)
650 NEXT I
655 PRINT D$;"CLOSE"
660 FOR M = 1 TO 7
670 CNUMB = SCHD(M)
675 PT = CP
680 GOSUB 1430
685 CREC(M) = REC
690 NEXT M
695 GOSUB 1500
700 GOSUB 1800
710 RETURN
1070 REM *****
1080 REM SUBROUTINE RETREIVE ALL STUDENT SCHEDULES
1085 REM *****
1110 PT = SP
1115 GOSUB 1220
1120 FOR T = 1 TO NST
1125 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
1130 PRINT D$;"READ STUDENT-SCHEDULE,R";ODR(T);" ,B0"
1135 INPUT SNUMB
1140 FOR M = 1 TO 7
1145 BYT = 6 + (M - 1) * 14
1150 PRINT D$;"READ STUDENT-SCHEDULE,R";ODR(T);" ,B";BYT
1155 INPUT SCHD(M)
1160 NEXT M
1165 PRINT D$;"CLOSE STUDENT-SCHEDULE"
1170 FOR M = 1 TO 7
1175 CNUMB = SCHD(M)
1180 PT = CP
1185 GOSUB 1425
1190 CREC(M) = REC
1195 NEXT M
1200 GOSUB 1495
1205 GOSUB 1555
1210 NEXT T

```

```

1215 RETURN
1220 REM *****
1225 REM SUBROUTINE IN-ORDER
1230 REM *****
1232 LN = 0
1235 LT = 0
1240 IF ANS = 3 AND SNDX(PT,1) < > 0 THEN GOTO 1310
1242 IF ANS = 4 AND TINDX(PT,1) < > 0 THEN GOTO 1310
1245 LN = LN + 1
1250 IF ANS = 3 THEN ODR(LN) = SNDX(PT,4)
1252 IF ANS = 4 THEN ODR(LN) = TINDX(PT,4)
1255 D = 1
1260 REM PUSH STACK
1265 GOSUB 1340
1270 D = 0
1275 IF ANS = 3 THEN PT = SNDX(PT,3)
1277 IF ANS = 4 THEN PT = TINDX(PT,3)
1280 IF PT < > 0 THEN GOTO 1240
1285 REM POP STACK
1290 GOSUB 1375
1295 IF D = 3 THEN GOTO 1335
1300 IF D = 1 THEN GOTO 1285
1305 GOTO 1245
1310 D = 2
1315 GOSUB 1340
1320 D = 0
1325 IF ANS = 3 THEN PT = SNDX(PT,1)
1327 IF ANS = 4 THEN PT = TINDX(PT,1)
1330 GOTO 1240
1335 RETURN
1340 REM *****
1345 REM SUBROUTINE PUSH
1350 REM *****
1355 LT = LT + 1
1360 STACK(LT,1) = PT
1365 STACK(LT,2) = D
1370 RETURN
1375 REM *****
1380 REM SUBROUTINE POP
1385 REM *****
1390 IF LT = 0 THEN GOTO 1415
1395 PT = STACK(LT,1)
1400 D = STACK(LT,2)
1405 LT = LT - 1
1410 GOTO 1420
1415 D = 3
1420 RETURN

```

```

1426 REM
1425 REM *****
1430 REM SUBROUTINE TO FIND CLASS RECORD
1435 REM *****
1440 IF PT = 0 THEN GOTO 1470
1445 IF CNUMB < CINDX(PT,2) THEN IK = 1
1450 IF CNUMB > CINDX(PT,2) THEN IK = 3
1455 IF CNUMB = CINDX(PT,2) THEN GOTO 1485
1460 PT = CINDX(PT,IK)
1465 GOTO 1440
1470 REM COURSE IS NOT IN DATA BASE
1475 FLAG = 1
1480 GOTO 1490
1485 REC = CINDX(PT,4)
1490 RETURN
1495 REM *****
1500 REM SUBROUTINE TO READ TEACHER AND ROOM NUMBER
1505 REM *****
1510 PRINT D$; "OPEN CLASS-INFO,L225"
1515 FOR I = 1 TO 7
1520 PRINT D$; "READ CLASS-INFO,R";CREC(I); ",B4"
1525 INPUT TCH(I)
1530 PRINT D$; "READ CLASS-INFO,R";CREC(I); ",B9"
1535 INPUT RM(I)
1540 NEXT I
1545 PRINT D$; "CLOSE CLASS-INFO"
1550 RETURN
1555 REM *****
1560 REM PRINT SUBROUTINE
1565 REM *****
1570 HOME : PRINT : PRINT
1575 PRINT "SCHEDULE FOR STUDENT ";SNUMB
1580 PRINT : PRINT
1585 PRINT "PERIOD"; TAB( 13) "COURSE";
1586 PRINT TAB( 24) "TEACHER"; TAB( 35) "ROOM"
1590 FOR I = 1 TO 7
1595 PRINT
1600 PRINT TAB( 3) I; TAB( 15) SCHD(I);
1605 PRINT TAB( 26) TCH(I); TAB( 36) RM(I)
1605 NEXT I
1610 PRINT : PRINT : INPUT "PRESS RETURN TO CONTINUE";ZZ$
1615 RETURN
1620 REM *****
1625 REM SUBROUTINE RETRIEVE ALL TEACHER SCHEDULES
1630 REM *****
1635 IF NTEA = 0 THEN GOTO 1770
1660 PT = TP

```



```

1665 GOSUB 1220
1670 FOR T = 1 TO NTEA
1675 PRINT D$;"OPEN TEACHER-SCHEDULE,L31"
1680 PRINT D$;"READ TEACHER-SCHEDULE,R";ODR(T);",B0"
1685 INPUT TNUMB
1690 FOR M = 1 TO 7
1695 BYT = 3 + (M - 1) * 4
1700 PRINT D$;"READ TEACHER-SCHEDULE,R";ODR(T);",B";BYT
1705 INPUT SCHD(M)
1710 NEXT M
1715 PRINT D$;"CLOSE TEACHER-SCHEDULE"
1720 FOR M = 1 TO 7
1725 CNUMB = SCHD(M)
1730 PT = CP
1735 GOSUB 1425
1740 CRC(M) = REC
1745 NEXT M
1750 GOSUB 1495
1755 GOSUB 1800
1760 NEXT T
1765 GOTO 1795
1770 REM TEACHER-INDEX HAS NOT BEEN BUILD
1775 HOME : PRINT : PRINT : PRINT
1780 PRINT "TEACHER-INDEX HAS NOT BEEN BUILD"
1785 PRINT : PRINT
1790 PRINT :PRINT:INPUT "PRESS RETURN TO CONTINUE ";ZZ$
1795 RETURN
1800 REM *****
1805 REM SUBROUTINE PRINT TEACHER SCHEDULE
1806 REM *****
1809 HOME : PRINT : PRINT
1810 PRINT "SCHEUDLE FOR TEACHER ";TNUMB
1815 PRINT : PRINT
1820 PRINT "PERIOD"; TAB( 15) "CLASS"; TAB( 25) "ROOM"
1825 FOR I = 1 TO 7
1830 PRINT
1835 PRINT TAB( 3)I; TAB( 16) SCHD(I); TAB( 26) RM(I)
1840 NEXT I
1845 PRINT : PRINT : INPUT "PRESS RETURN TO CONTINUE ";ZZ$
1850 RETURN

```

APPENDIX E

PROGRAM UPDATE FILES

```

100 REM PROGRAM TO UPDATE RECORDS
110 REM
120 DIM SNDX(500,4)
130 DIM TINDX(50,4)
140 DIM CINDX(150,4)
150 DIM SCHD(7)
160 DIM CH(7,2)
170 D$ = CHR$(4)
180 GOSUB 450
190 GOSUB 570
200 GOSUB 700
210 HOME : PRINT : PRINT
220 PRINT "THIS PROGRAM UPDATES:"
230 PRINT : PRINT
240 PRINT "ENTER THE NUMBER OF YOUR CHOICE"
250 PRINT
260 PRINT "1.....STUDENT-SCHEDULE"
270 PRINT
280 PRINT "2.....TEACHER-SCHEDULE"
290 PRINT
300 PRINT "3.....STUDENT GRADES"
310 PRINT
320 PRINT "4.....ROOM NUMBERS"
330 PRINT
340 PRINT "5.....ENTER 5 TO END PROGRAM"
350 PRINT : PRINT : INPUT ANS
360 IF ANS = 1 THEN GOSUB 830
370 IF ANS = 2 THEN GOSUB 1810
380 IF ANS = 3 THEN GOSUB 2960
390 IF ANS = 4 THEN GOSUB 3680
400 IF ANS = 5 THEN GOTO 430
410 IF NCLASS < > CLASS0 THEN GOSUB 2840
420 GOTO 210
430 REM EXIT PROGRAM AND RUN MENU PROGRAM
440 PRINT D$;"RUN MENU,D1"
450 REM *****
460 REM SUBROUTINE TO READ ST-INDEX"
470 REM *****

```

```

480 PRINT D$;"OPEN STUDENT-INDEX,D2"
490 PRINT D$;"READ STUDENT-INDEX"
500 INPUT NST: INPUT SP: INPUT OBAL
510 FOR I = 1 TO NSTU
520 FOR J = 1 TO 4
530 INPUT SNDX(I,J)
540 NEXT J: NEXT I
550 PRINT D$;"CLOSE ST-INDEX"
560 RETURN
570 REM *****
580 REM SUBROUTINE TO READ CLASS-INDEX"
590 REM *****
600 PRINT D$;"OPEN CLASS-INDEX"
610 PRINT D$;"READ CLASS-INDEX"
620 INPUT NCLASS: INPUT CP: INPUT OCBAL
630 CLASS0 = NCLASS
640 FOR I = 1 TO NCLASS
650 FOR J = 1 TO 4
660 INPUT CINDX(I,J)
670 NEXT J: NEXT I
680 PRINT D$;"CLOSE CLASS-INDEX"
690 RETURN
700 REM *****
710 REM SUBROUTINE TO READ TEACHER-INDEX"
720 REM *****
730 PRINT D$;"OPEN TEACHER-INDEX"
740 PRINT D$;"READ TEACHER-INDEX"
750 INPUT NTEA: INPUT TP: INPUT OBAL
760 IF NTEA = 0 THEN GOTO 810
770 FOR I = 1 TO NTEA
780 FOR J = 1 TO 4
790 INPUT TINDX(I,J)
800 NEXT J: NEXT I
810 PRINT D$;"CLOSE TEACHER-INDEX"
820 RETURN
830 REM *****
840 REM SUBROUTINE STUDENT-SCHEDULE
850 REM *****
860 HOME : PRINT : PRINT
870 PRINT "ENTER STUDENT NUMBER"
880 PRINT : PRINT : INPUT NUMB
890 REM FIND STUDENT RECORD NUMBER
900 PT = SP
910 IF PT = 0 THEN GOTO 970
920 IF NUMB < SNDX(PT,2) THEN IK = 1
930 IF NUMB > SNDX(PT,2) THEN IK = 3

```

```

940 IF NUMB = SNDX(PT,2) THEN GOTO 1000
950 PT = SNDX(PT,IK)
960 GOTO 910
970 REM STUDENT IS NOT IN DATA BASE
980 GOSUB 2490
990 GOTO 1800
1000 REC = SNDX(PT,4)
1010 TREC = REC
1020 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
1030 FOR I = 1 TO 7
1040 BYT = 6 + (I - 1) * 14
1050 PRINT D$;"READ STUDENT-SCHEDULE,R";REC;"B";BYT
1060 INPUT SCHD(I)
1070 NEXT I
1080 PRINT D$;"CLOSE STUDENT-SCHEDULE"
1090 HOME : PRINT : PRINT
1100 PRINT "STUDENT ";NUMB;" CURRENT SCHEDULE IS:" : PRINT
1110 PRINT "PERIOD CLASS"
1120 FOR I = 1 TO 7
1130 PRINT
1140 PRINT " ";I;" ";SCHD(I)
1150 NEXT I
1160 PRINT : PRINT
1170 INPUT "ENTER PERIOD, NEW COURSE ";PN,COURSE
1180 CNUMB = SCHD(PN)
1190 FLAG = 1
1200 GOSUB 2630
1210 IF FLAG = 3 THEN GOTO 1720
1220 FLAG = 0
1230 SCHD(PN) = COURSE
1240 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
1250 BYT = 6 + (PN - 1) * 14
1260 PRINT D$;"WRITE STUDENT-SCHEDULE,R";TREC;"B";BYT
1270 PRINT SCHD(PN)
1280 REM UPDATE CLASS-INFO RELATION
1290 PRINT D$;"CLOSE STUDENT-SCHEDULE"
1300 REM UPDATE OLD CLASS LIST
1310 PRINT D$;"OPEN CLASS-INFO,L225"
1320 REM DECREMENT CSIZE
1330 PRINT D$;"READ CLASS-INFO,R";REC;"B12"
1340 INPUT CSIZE
1350 CSIZE = CSIZE - 1
1360 PRINT D$;"WRITE CLASS-INFO,R";REC;"B12"
1370 PRINT CSIZE
1380 FOR I = 1 TO (CSIZE + 1)
1390 BYT = 15 + (I - 1) * 6
1400 PRINT D$;"READ CLASS-INFO,R";REC;"B";BYT

```

```

1410 INPUT LST(I)
1420 NEXT I
1430 REM FIND AND DELETE STUDENT
1440 FOR J = 1 TO (CSIZE + 1)
1450 IF NUMB = LST(J) THEN GOTO 1470
1460 NEXT J
1470 IF J = (CSIZE + 1) THEN GOTO 1540
1480 FOR I = J TO CSIZE
1490 BYT = 15 + (I - 1) * 6
1500 PRINT D$; "WRITE CLASS-INFO,R";REC; ",B";BYT
1510 PRINT LST(I + 1)
1520 NEXT I
1530 GOTO 1570
1540 BYT = 15 + CSIZE * 6
1550 PRINT D$; "WRITE CLASS-INFO,R";REC; ",B";BYT
1560 PRINT 0
1570 REM ADD STUDENT TO NEW CLASS LIST
1580 CNUMB = SCHD(PN)
1590 FLAG = 0
1600 GOSUB 2630
1610 PRINT D$; "READ CLASS-INFO,R";REC; ",B12"
1620 INPUT CSIZE
1630 CSIZE = CSIZE + 1
1640 PRINT D$; "WRITE CLASS-INFO,R";REC; ",B12"
1650 PRINT CSIZE
1660 BYT = 15 + (CSIZE - 1) * 6
1670 PRINT D$; "WRITE CLASS-INFO,R";REC; ",B";BYT
1680 PRINT NUMB
1690 PRINT D$; "CLOSE CLASS-INFO"
1700 GOTO 1740
1710 PRINT D$; "CLOSE CLASS-INFO"
1720 PRINT : PRINT "STUDENT ";NUMB;
1725 PRINT " IS NOT ENROLLED IN COURSE ";COURSE;
1726 PRINT " IN PERIOD ";PN
1730 FLAG = 0
1740 HOME : PRINT : PRINT : PRINT
1750 PRINT "DO YOU WANT TO MAKE ANY MORE CHANGES"
1760 PRINT " FOR THIS STUDENT "
1770 INPUT "... (Y/N) ..";A$
1780 IF A$ = "Y" THEN GOTO 1090
1790 IF A$ < > "N" THEN GOTO 1750
1800 RETURN
1810 REM *****
1820 REM SUBROUTINE TEACHER-SCHEDULE
1830 REM *****
1840 HOME : PRINT : PRINT
1850 PRINT "ENTER TEACHER NUMBER"

```

```

1860 PRINT : PRINT : INPUT NUMB
1870 PT = TP
1880 IF PT = 0 THEN GOTO 1940
1890 IF NUMB < TINDX(PT,2) THEN IK = 1
1900 IF NUMB > TINDX(PT,2) THEN IK = 3
1910 IF NUMB = TINDX(PT,2) THEN GOTO 1970
1920 PT = TINDX(PT,IK)
1930 GOTO 1880
1940 REM TEACHER IS NOT IN DATA BASE
1950 GOSUB 2490
1960 GOTO 2480
1970 TREC = TINDX(PT,4)
1980 PRINT D$;"OPEN TEACHER-SCHEDULE,L31"
1990 FOR I = 1 TO 7
2000 BYT = 3 + (I - 1) * 4
2010 PRINT D$;"READ TEACHER-SCHEDULE,R";TREC;"B";BYT
2020 INPUT SCHD(I)
2030 NEXT I
2040 PRINT D$;"CLOSE TEACHER-SCHEDULE"
2050 HOME : PRINT : PRINT : PRINT
2060 PRINT "TEACHER ";NUMB;" CURRENT SCHEDULE IS: "
2070 PRINT : PRINT "PERIOD CLASS"
2080 FOR I = 1 TO 7
2090 PRINT
2100 PRINT I;" ";SCHD(I)
2110 NEXT I
2120 PRINT : PRINT
2130 INPUT "ENTER PERIOD,NEW COURSE";PN,COURSE
2140 IF SCHD(PN) < > COURSE THEN GOTO 2180
2150 PRINT "TEACHER ALREADY ASSIGNED TO CLASS ";CN
2155 PRINT " IN PERIOD ";PN
2160 INPUT "PRESS RETURN TO CONTINUE";ZZ$
2170 GOTO 2235
2180 FLAG = 1
2185 CNUMB = SCHD(PN)
2190 GOSUB 2630
2200 IF FLAG = 3 THEN GOTO 2420
2210 SCHD(PN) = COURSE
2220 PRINT D$;"OPEN TEACHER-SCHEDULE,L31"
2230 BYT = 3 + (PN - 1) * 4
2240 PRINT D$;"WRITE TEACHER-SCHEDULE,R";TREC;"B";BYT
2250 PRINT SCHD(PN)
2260 PRINT D$;"CLOSE TEACHER-SCHEDULE"
2270 ROLD = REC
2280 CNUMB = SCHD(PN)
2290 FLAG = 0
2300 GOSUB 2630

```

```

2310 RNW = REC
2320 PRINT D$; "OPEN CLASS-INFO,L225"
2330 PRINT D$; "READ CLASS-INFO,R";ROLD; ",B4"
2340 INPUT TEMP
2350 IF TEMP < > NUMB GOTO 2380
2360 PRINT D$; "WRITE CLASS-INFO,R";ROLD; ",B4"
2370 PRINT 0
2380 PRINT D$; "WRITE CLASS-INFO,R";RNW; ",B4"
2390 PRINT NUMB
2400 PRINT D$; "CLOSE CLASS-INFO"
2410 GOTO 2430
2420 PRINT : PRINT "TEACHER ";NUMB;
2415 PRINT " IS NOT IN COURSE ";COURSE;" IN PERIOD ";PN
2430 PRINT : PRINT : "DO YOU WANT TO MAKE ANY MORE CHANGES"
2440 PRINT : PRINT "      FOR THIS TEACHER"
2450 PRINT : INPUT "....(Y/N) ....";A$
2460 IF A$ = "Y" THEN GOTO 2050
2470 IF A$ < > "N" THEN GOTO 2430
2480 RETURN
2490 REM *****
2500 REM SUBROUTINE MISSING-PERSON
2510 REM *****
2520 HOME : PRINT : PRINT : PRINT : PRINT : PRINT
2530 IF ANS = 1 THEN PRINT "STUDENT";
2540 IF ANS = 2 THEN PRINT "TEACHER";
2550 PRINT " ";NUMB;" IS NOT IN THE INDEX"
2560 PRINT : PRINT : PRINT :
2570 PRINT "USE THE BUILD PROGRAM TO ENTER"
2580 PRINT : PRINT
2590 PRINT "TEACHERS AND STUDENTS INTO THE DATA BASE"
2600 PRINT : PRINT : PRINT : PRINT
2610 INPUT "PRESS RETURN TO CONTINUE";ZZ$
2620 RETURN
2630 REM *****
2640 REM SUBROUTINE TO FIND CLASS RECORD NUMBER
2650 REM *****
2660 PT = CP
2670 IF PT = 0 THEN GOTO 2750
2680 IF CNUMB < CINDX(PT,2) THEN IK = 1
2690 IF CNUMB > CINDX(PT,2) THEN IK = 3
2700 IF CNUMB = CINDX(PT,2) THEN GOTO 2730
2710 PT = CINDX(PT,IK)
2720 GOTO 2670
2730 REC = CINDX(PT,4)
2740 GOTO 2830
2750 IF FLAG = 1 THEN GOTO 2820
2760 NCLASS = NCLASS + 1

```

```

2770 CINDX(LAST,IK) = NCLASS
2780 CINDX(NCLASS,2) = SCHD(PN)
2790 CINDX(NCLASS,4) = NCLASS - 1
2800 REC = NCLASS - 1
2810 GOTO 2830
2820 FLAG = 3
2830 RETURN
2840 REM *****
2850 REM SUBROUTINE WRITE CLASS-INDEX
2860 REM *****
2870 PRINT D$;"OPEN CLASS-INDEX"
2880 PRINT D$;"OPEN CLASS-INDEX"
2890 PRINT NCLASS: PRINT CP: PRINT OCBAL
2900 FOR I = 1 TO NCLASS
2910 FOR J = 1 TO 4
2920 PRINT CINDX(I,J)
2930 NEXT J: NEXT I
2940 PRINT D$;"CLOSE CLASS-INDEX"
2950 RETURN
2960 REM *****
2970 REM SUBROUTINE UPDATE-GRADES
2980 REM *****
2990 HOME : PRINT : PRINT : PRINT
3000 PRINT "1.....TO UPDATE MID-SEMESTER GRADE"
3010 PRINT : PRINT
3020 PRINT "2.....TO UPDATE FINIAL-GRADE"
3030 PRINT : PRINT : PRINT
3040 PRINT "ENTER THE NUMBER OF YOUR CHOICE"
3050 INPUT BANS
3060 IF BANS = 1 THEN GOSUB 3560
3070 IF BANS = 2 THEN GOSUB 3620
3080 IF BANS < > 1 AND BANS < > 2 THEN GOTO 2990
3090 HOME : PRINT : PRINT : PRINT
3100 INPUT "ENTER STUDENT NUMBER ";NUMB
3110 PT = SP
3120 IF PT = 0 THEN GOTO 3180
3130 IF NUMB < SNDX(PT,2) THEN IK = 1
3140 IF NUMB > SNDX(PT,2) THEN IK = 3
3150 IF NUMB = SNDX(PT,2) THEN GOTO 3210
3160 PT = SNDX(PT,IK)
3170 GOTO 3120
3180 PRINT : PRINT "STUDENT ";NUMB;
3185 PRINT " IS NOT IN THE DATA BASE"
3190 PRINT : INPUT "PRESS RETURN TO CONTINUE";ZZ$
3200 GOTO 3550
3210 REC = SNDX(PT,4)
3220 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"

```



```

3230 FOR I = 1 TO 7
3240 BYT = 6 + (I - 1) * 14
3250 PRINT D$; "READ STUDENT--SCHEDULE,R"; REC; ",B"; BYT
3260 INPUT SCHD(I)
3270 BYT = B1 + (I - 1) * 14
3280 PRINT D$; "READ STUDENT--SCHEDULE,R"; REC; ",B"; BYT
3290 INPUT GDE(I)
3300 BYT = B2 + (I - 1) * 14
3310 PRINT D$; "READ STUDENT--SCHEDULE,R"; REC; ",B"; BYT
3320 INPUT COM(I)
3330 NEXT I
3340 PRINT D$; "CLOSE STUDENT-SCHEDULE"
3350 HOME : PRINT : PRINT
3360 PRINT "CURRENT GRADES FOR STUDENT "; NUMB; " ARE:"
3370 PRINT : PRINT
3380 PRINT "PERIOD"; TAB( 11) "CLASS"; TAB( 20) "GRADE";
3385 PRINT TAB( 30) "COMMENT"
3390 FOR I = 1 TO 7
3400 PRINT
3410 PRINT TAB( 3) I; TAB( 12) SCHD(I); TAB( 22) GDE(I);
3405 PRINT TAB( 33) COM(I)
3420 NEXT I
3430 PRINT
3440 PRINT "ENTER PERIOD, NEW GRADE, NEW COMMENT"
3450 INPUT PN, GD, CM
3460 IF PN > 7 OR PN < 1 THEN GOTO 3440
3470 PRINT D$; "OPEN STUDENT-SCHEDULE,L104"
3480 BYT = B1 + (PN - 1) * 14
3490 PRINT D$; "WRITE STUDENT-SCHEDULE,R"; REC; ",B"; BYT
3500 PRINT GD
3510 BYT = B2 + (PN - 1) * 14
3520 PRINT D$; "WRITE STUDENT-SCHEDULE,R"; REC; ",B"; BYT
3530 PRINT CM
3540 PRINT D$; "CLOSE STUDENT-SCHEDULE"
3550 RETURN
3560 REM *****
3570 REM SUBROUTINE BYTE POINTERS FOR MID-SEMESTER
3580 REM *****
3590 B1 = 10
3600 B2 = 13
3610 RETURN
3620 REM *****
3630 REM SUBROUTINE BYTE-POINTERS-MID-SEMESTER
3640 REM *****
3650 B1 = 15
3660 B2 = 18
3670 RETURN

```

```

3675 REM
3680 REM *****
3690 REM SUBROUTINE CHANGE-ROOM
3700 REM *****
3710 HOME : PRINT : PRINT : PRINT
3720 INPUT "ENTER THE CLASS NUMBER ";CNUMB
3730 PT = CP
3740 IF PT = 0 THEN GOTO 3800
3750 IF CNUMB < CINDX(PT,2) THEN IK = 1
3760 IF CNUMB > CINDX(PT,2) THEN IK = 3
3770 IF CNUMB = CINDX(PT,2) THEN GOTO 3830
3780 PT = CINDX(PT,IK)
3790 GOTO 3740
3800 PRINT : PRINT : PRINT "COURSE ";CNUMB;
3805 PRINT "IS NOT IN THE DATA BASE"
3810 INPUT "PRESS RETURN TO CONTINUE ";ZZ$
3820 GOTO 3970
3830 REC = CINDX(PT,4)
3840 PRINT D$;"OPEN CLASS-INFO,L225"
3850 PRINT D$;"READ CLASS-INFO,R";REC;" ,B9"
3860 INPUT RN
3870 PRINT D$;"CLOSE CLASS-INFO"
3880 HOME : PRINT : PRINT : PRINT
3890 PRINT "CLASS ";CNUMB;" IS CURRENTLY SCHEDULED IN"
3900 PRINT : PRINT " ROOM ";RN
3910 PRINT : PRINT : PRINT : PRINT
3920 INPUT "ENTER NEW ROOM NUMBER ";NRM
3930 PRINT D$;"OPEN CLASS-INFO,L225"
3940 PRINT D$;"WRITE CLASS-INFO,R";REC;" ,B9"
3950 PRINT NRM
3960 PRINT D$;"CLOSE CLASS-INFO"
3970 RETURN

```

APPENDIX F

PROGRAM ENTER GRADES

```

100 REM
110 REM PROGRAM TO ENTER GRADES
120 REM
130 DIM TINDX(50,4)
140 DIM CINDX(150,4)
150 DIM SCHD(7)
160 DIM GDE(35)
170 DIM COM(35)
180 DIM SREC(35)
190 D$ = CHR$(4)
200 GOSUB 740
210 GOSUB 870
220 GOSUB 1000
230 HOME : PRINT : PRINT
240 PRINT : PRINT : PRINT : PRINT
250 PRINT "1.....FOR MID-SEMESTER GRADES"
260 PRINT : PRINT
270 PRINT "2.....OR FINIAL GRADES"
280 PRINT : PRINT
290 INPUT "ENTER NUMBER OF YOUR CHOICE";ANS
300 IF ANS = 1 THEN GOSUB 620
310 IF ANS = 2 THEN GOSUB 680
320 IF ANS < > 1 AND ANS < > 2 THEN GOTO 230
330 HOME : PRINT : PRINT : PRINT
340 PRINT "ENTER THE TEACHER NUMBER"
350 PRINT : PRINT : PRINT
360 PRINT "OR"
370 PRINT : PRINT : PRINT
380 PRINT "ENTER ..-9..TO TERMINATE PROGRAM"
390 PRINT : PRINT
400 INPUT NUMB
410 IF NUMB = - 9 THEN GOTO 600
420 GOSUB 1130
430 IF FTEAC = 1 THEN GOTO 530
440 GOSUB 1280
450 FOR KK = 1 TO 7
460 NUMB = SCHD(KK)
470 IF NUMB = 999 GOTO 510

```

```

480 GOSUB 1390
490 GOSUB 1530
500 GOSUB 1920
510 NEXT KK
520 GOTO 310
530 REM GIVEN TEACHER IS NOT IN DATA BASE
540 PRINT : PRINT : PRINT
550 PRINT "TEACHER ";NUMB;" IS NOT IN THE DATA BASE"
560 PRINT : PRINT : PRINT
570 INPUT "PRESS THE RETURN KEY TO CONTINUE";ZZ$
580 FTEACH = 0
590 GOTO 310
600 REM EXIT PROGRAM AND RUN MENU
610 PRINT D$;"RUN MENU,D1"
620 REM *****
630 REM SUBROUTINE TO SET BYTE POINTERS
640 REM *****
650 B1 = 10
660 B2 = 13
670 RETURN
680 REM *****
690 REM SUBROUTINE TO SET BYTE POINTERS FOR FINIAL GRADES
700 REM *****
710 B1 = 15
720 B2 = 18
730 RETURN
740 REM *****
750 REM SUBROUTINE TO READ TEACHER-INDEX"
760 REM *****
770 PRINT D$;"OPEN TEACHER-INDEX,D2"
780 PRINT D$;"READ TEACHER-INDEX"
790 INPUT NTES: INPUT TP: INPUT OBAL
800 IF NTEA = 0 THEN GOTO 850
810 FOR I = 1 TO NTEA
820 FOR J = 1 TO 4
830 INPUT TINDX(I,J)
840 NEXT J: NEXT I
850 PRINT D$;"CLOSE TEACHER-INDEX"
860 RETURN
870 REM *****
880 REM SUBROUTINE TO READ CLASS-INDEX"
890 REM *****
900 PRINT D$;"OPEN CLASS-INDEX"
910 PRINT D$;"READ CLASS-INDEX"
920 INPUT NCLASS: INPUT CP: INPUT OCBAL
930 IF NCLASS = 0 THEN GOTO 980

```

```

940 FOR I = 1 TO NCLASS
950 FOR J = 1 TO 4
960 INPUT CINDX(I,J)
970 NEXT J: NEXT I
980 PRINT D$;"CLOSE CLASS-INDEX"
990 RETURN
1000 REM *****
1010 REM SUBROUTINE TO READ STUDENT-INDEX"
1020 REM *****
1030 PRINT D$;"OPEN STUDENT-INDEX"
1040 PRINT D$;"READ STUDENT-INDEX"
1050 INPUT NST: INPUT SP: INPUT OBAL
1060 IF NSTU = 0 THEN GOTO 1110
1070 FOR I = 1 TO NSTU
1080 FOR J = 1 TO 4
1090 INPUT SNDX(I,J)
1100 NEXT J: NEXT I
1110 PRINT D$;"CLOSE ST-INDEX"
1120 RETURN
1130 REM *****
1140 REM SUBROUTINE TO FIND TEACHER RECORD"
1150 REM *****
1160 PT = TP
1170 IF PT = 0 THEN GOTO 1230
1180 IF NUMB < TINDX(PT,2) THEN IK = 1
1190 IF NUMB > TINDX(PT,2) THEN IK = 3
1200 IF NUMB = TINDX(PT,2) THEN GOTO 1260
1210 PT = TINDX(PT,IK)
1220 GOTO 1170
1230 REM TEACHER IS NOT IN DATA BASE
1240 FTEAC = 1
1250 RETURN
1260 REC = TINDX(PT,4)
1270 RETURN
1280 REM *****
1290 REM SUBROUTINE TO READ TEACHER SCHEDULE
1300 REM *****
1310 PRINT D$;"OPEN TEACHER-SCHEDULE,L31"
1320 FOR J = 1 TO 7
1330 BYT = 3 + (J - 1) * 4
1340 PRINT D$;"READ TEACHER-SCHEDULE,R";REC;" ,B";BYT
1350 INPUT SCHD(J)
1360 NEXT J
1370 PRINT D$;"CLOSE TEACHER-SCHEDULE"
1380 RETURN
1382 REM
1385 REM

```

```

1387 REM
1390 REM *****
1400 REM SUBROUTINE TO FIND CLASS RECORD
1410 REM *****
1420 PT = CP
1430 IF PT = 0 THEN GOTO 1490
1440 IF NUMB < CINDX(PT,2) THEN IK = 1
1450 IF NUMB > CINDX(PT,2) THEN IK = 3
1460 IF NUMB = CINDX(PT,2) THEN 1510
1470 PT = CINDX(PT,IK)
1480 GOTO 1430
1490 REM CLASS IS NOT IN DATA BASE
1500 FLAG = 1
1510 REC = CINDX(PT,4)
1520 RETURN
1530 REM *****
1540 REM SUBROUTINE TO READ CLASS-LIST
1550 REM *****
1560 PRINT D$;"OPEN CLASS-INFO,L225"
1570 PRINT D$;"READ CLASS-INFO,R";REC;"",B0"
1580 INPUT INFO(1)
1590 PRINT D$;"READ CLASS-INFO,R";REC;"",B4"
1600 INPUT INFO(2)
1610 PRINT D$;"READ CLASS-INFO,R";REC;"",B9"
1620 INPUT INFO(4)
1630 PRINT D$;"READ CLASS-INFO,R";REC;"",B7"
1635 INPUT INFO(3)
1640 PRINT D$;"READ CLASS-INFO,R";REC;"",B12"
1650 INPUT INFO(5)
1660 IF INFO(5) = 0 THEN 1740
1670 FOR I = 1 TO INFO(5)
1680 BYT = 15 + (I - 1) * 6
1690 PRINT D$;"READ CLASS-INFO,R";REC;"",B";BYT
1700 INPUT LST(I)
1710 NEXT I
1720 GOTO 1740
1730 FLAG = 1
1740 PRINT D$;"CLOSE CLASS-INFO"
1750 RETURN
1760 REM *****
1770 REM SUBROUTINE TO FIND STU-SCH RECORDS
1780 REM *****
1790 FOR I = 1 TO INFO(5)
1800 PT = SP
1810 IF PT = 0 THEN GOTO 1890
1820 IF LST(I) < SNDX(PT,2) THEN IK = 1
1830 IF LST(I) > SNDX(PT,2) THEN IK = 3

```

```

1840 IF LST(I) = SNDX(PT,2) THEN 1870
1850 PT = SNDX(PT,IK)
1860 GOTO 1810
1870 SREC(I) = SNDX(PT,4)
1880 GOTO 1900
1890 SREC(I) = - 9
1900 NEXT I
1910 RETURN
1920 REM *****
1930 REM SUBROUTINE TO INPUT AND PRINT GRADES
1940 REM *****
1950 HOME : PRINT : PRINT
1960 PRINT "CLASS TEACHER PERIOD RM# CSIZE"
1970 PRINT
1980 FOR I = 1 TO 5
1990 PRINT INFO(I); " ";
2000 NEXT I
2010 PRINT : PRINT
2020 IF INFO(5) = 0 THEN GOTO 2320
2030 PRINT "ENTER GRADE,COMMENT"
2040 PRINT
2050 PRINT "STUDENT GRADE,COMMENT"
2060 PRINT
2070 FOR J = 1 TO INFO(5)
2080 PRINT LST(J); " ";
2090 INPUT GDE(J),COM(J)
2100 PRINT
2110 NEXT J
2120 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
2130 REM WRITE TO STUDENT-SCHEDULE FILE
2140 FOR I = 1 TO INFO(5)
2150 PT = SP
2160 IF PT = 0 THEN GOTO 2230
2170 IF LST(I) < SNDX(PT,2) THEN IK = 1
2180 IF LST(I) = SNDX(PT,2) THEN GOTO 2220
2190 IF LST(I) > SNDX(PT,2) THEN IK = 3
2200 PT = SNDX(PT,IK)
2210 GOTO 2160
2220 SREC(I) = SNDX(PT,4)
2230 NEXT I
2240 FOR I = 1 TO INFO(5)
2250 BYT = B1 + (INFO(2)-1) * 14
2260 PRINT D$;"WRITE STUDENT-SCHEDULE,R";SREC(I);",B";BYT"
2270 PRINT GDE(I)
2280 BYT = B2 + (INFO(2) - 1) * 14
2290 PRINT D$;"WRITE STUDENT-SCHEDULE,R";SREC(I);",B";BYT"
2300 PRINT COM(I)

```

```
2310 NEXT I
2320 PRINT D$;"CLOSE STUDENT-SCHEDULE "
2330 INPUT "PRESS RETURN FOR THE NEXT CLASS";ZZ$
2340 RETURN
```


APPENDIX G

PROGRAM PRINT GRADES

```

100 REM
110 REM PROGRAM PRINT-GRADES
120 REM
130 REM
140 DS = CHR$ (4)
150 DIM SNDX(500,4)
160 DIM CINDX(150,4)
170 DIM SCHD(7)
180 DIM TCH(7)
190 DIM GMT(7),GF(7)
200 DIM CMT(7),CF(7)
210 DIM CREC(7)
220 DIM STACK(500,2)
230 DIM ODR(500)
240 FLAG = 0
250 GOSUB 730
260 GOSUB 880
270 IF FLAG = 1 THEN GOTO 580
280 HOME : PRINT : PRINT
290 PRINT "THIS PROGRAM PRINTS THE: "
300 PRINT : PRINT : PRINT
310 PRINT "1.....MID-SEMESTER GRADES."
320 PRINT : PRINT : PRINT
330 PRINT "2.....FINIAL GRADES."
340 PRINT : PRINT : PRINT
350 PRINT "3.....TO END PROGRAM"
360 PRINT : PRINT : PRINT
370 INPUT "ENTER THE NUMBER OF YOUR CHOICE: ";ANS
380 IF ANS = 3 THEN GOTO 630
390 IF ANS < > 1 AND ANS < > 2 THEN GOTO 370
400 GOSUB 650
410 HOME : PRINT : PRINT : PRINT
420 PRINT "DO YOU WANT GRADES FOR: "
430 PRINT : PRINT
440 PRINT "1.....ONE STUDENT"
450 PRINT : PRINT
460 PRINT "2.....ALL STUDENTS"
470 PRINT : PRINT

```

```

475 INPUT "ENTER THE NUMBER OF YOUR CHOICE: ";BANS
480 IF BANS = 1 THEN GOSUB 1030
490 IF BANS = 2 THEN GOSUB 1200
500 IF BANS < > 1 AND BANS < > 2 THEN GOTO 420
510 HOME : PRINT : PRINT
520 PRINT "DO YOU WANT TO RETRIEVE MORE"
530 PRINT : PRINT "      STUDENT'S GRADES (Y/N) "
540 PRINT : PRINT : INPUT C$
550 IF C$ = "Y" THEN GOTO 410
560 IF C$ = "N" THEN GOTO 280
570 GOTO 510
580 REM EXIT PROGRAM AND RUN MENU PROGRAM
590 HOME : PRINT : PRINT
600 PRINT "THE CLASS INDEX HAS NOT BEEN BUILT"
610 PRINT:PRINT:PRINT " USE BUILD PROGRAM TO ENTER CLASSES"
620 PRINT : PRINT : INPUT "PRESS RETURN TO CONTINUE ";ZZ$
630 REM EXIT PROGRAM AND RUN MENU
640 PRINT D$;"RUN MENU,D1"
650 REM *****
660 REM SUBROUTINE SET BYTE-POINTERS
670 REM *****
680 B1 = 10
690 B2 = 13
700 B3 = 15
710 B4 = 18
720 RETURN
730 REM *****
740 REM SUBROUTINE READ STUDENT-INDEX
750 REM *****
760 PRINT D$;"OPEN STUDENT-INDEX,D2"
770 PRINT D$;"READ STUDENT-INDEX"
780 INPUT NST: INPUT SP: INPUT OBAL
790 IF NST = 0 THEN GOTO 850
800 FOR I = 1 TO NST
810 FOR J = 1 TO 4
820 INPUT SNDX(I,J)
830 NEXT J: NEXT I
840 GOTO 860
850 FLAG = 1
860 PRINT D$;"CLOSE STUDENT-INDEX"
870 RETURN
880 REM *****
890 REM SUBROUTINE READ CLASS-INDEX"
900 REM *****
910 PRINT D$;"OPEN CLASS-INDEX"
920 PRINT D$;"READ CLASS-INDEX"

```

```

930 INPUT NCLASS: INPUT CP: INPUT OCBAL
940 IF NCLASS = 0 THEN GOTO 1000
950 FOR I = 1 TO NCLASS
960 FOR J = 1 TO 4
970 INPUT CINDX(I,J)
980 NEXT J: NEXT I
990 GOTO 1010
1000 FLAG = 1
1010 PRINT D$;"CLOSE CLASS-INDEX"
1020 RETURN
1030 REM *****
1040 REM SUBROUTINE PRINT-GRADES (ONE STUDENT)
1050 REM *****
1060 HOME : PRINT
1070 INPUT "ENTER THE STUDENT NUMBER ";NUMB
1080 FLAG = 0
1090 GOSUB 1310
1100 IF FLAG = 1 THEN GOTO 1150
1110 GOSUB 1450
1120 IF ANS = 1 THEN GOSUB 2340
1130 IF ANS = 2 THEN GOSUB 2480
1140 GOTO 1190
1150 HOME : PRINT : PRINT
1160 PRINT " STUDENT ";NUMB;" IS NOT IN THE DATA BASE"
1170 PRINT : PRINT : PRINT
1180 INPUT "PRESS RETURN TO CONTINUE F ";ZZ$
1190 RETURN
1200 REM *****
1210 REM SUBROUTINE PRINT-GRADES (ALL STUDENTS)
1220 REM *****
1230 GOSUB 1950
1240 FOR JJ = 1 TO NST
1250 REC = SNDX(ODR(JJ),4)
1260 GOSUB 1450
1270 IF ANS = 1 THEN GOSUB 2340
1280 IF ANS = 2 THEN GOSUB 2480
1290 NEXT JJ
1300 RETURN
1310 REM *****
1320 REM SUBROUTINE READ STUDENT-RECORD
1330 REM *****
1340 PT = SP
1350 IF PT = 0 THEN GOTO 1430
1360 IF NUMB < SNDX(PT,2) THEN IK = 1
1370 IF NUMB > SNDX(PT,2) THEN IK = 3
1380 IF NUMB = SNDX(PT,2) THEN GOTO 1410
1390 PT = SNDX(PT,IK)

```

```

1400 GOTO 1350
1410 REC = SNDX(PT,4)
1420 GOTO 1440
1430 FLAG = 1
1440 RETURN
1450 REM *****
1460 REM SUBROUTINE READ-GRADES
1470 REM *****
1480 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
1490 PRINT D$;"READ STUDENT-SCHEDULE,R";REC;"",B0"
1500 INPUT NUMB
1510 FOR I = 1 TO 7
1520 BYT = 6 + (I - 1) * 14
1530 PRINT D$;"READ STUDENT-SCHEDULE,R";REC;"",B";BYT
1540 INPUT SCHD(I)
1550 NEXT I
1560 PRINT D$;"CLOSE STUDENT-SCHEDULE"
1570 GOSUB 1720
1580 PRINT D$;"OPEN STUDENT-SCHEDULE,L104"
1590 FOR I = 1 TO 7
1595 BYT = B1 + (I-1) * 14
1600 PRINT D$;"READ STUDENT-SCHEDULE,R";REC;"",B";BYT
1610 INPUT GMT(I)
1615 BYT = B2 + (I-1) * 14
1620 PRINT D$;"READ STUDENT-SCHEDULE,R";REC;"",B";BYT
1630 INPUT CMT(I)
1640 IF ANS = 1 THEN GOTO 1690
1645 BYT = B3 + (I-1) * 14
1650 PRINT D$;"READ STUDENT-SCHEDULE,R";REC;"",B";BYT
1660 INPUT GF(I)
1665 BYT = B4 + (I-1) * 14
1670 PRINT D$;"READ STUDENT-SCHEDULE,R";REC;"",B";BYT
1680 INPUT CF(I)
1690 NEXT I
1700 PRINT D$;"CLOSE STUDENT-SCHEDULE"
1710 RETURN
1720 REM *****
1730 REM SUBROUTINE READ CLASS-DATA
1740 REM *****
1750 PRINT D$;"OPEN CLASS-INFO,L225"
1760 FOR I = 1 TO 7
1770 CNUMB = SCHD(I)
1780 PT = CP
1790 IF PT = 0 THEN GOTO 1910
1800 IF CNUMB < CINDX(PT,2) THEN IK = 1
1810 IF CNUMB > CINDX(PT,2) THEN IK = 3
1820 IF CNUMB = CINDX(PT,2) THEN GOTO 1850

```

```

1830 PT = CINDX (PT,IK)
1840 GOTO 1790
1850 CREC = CINDX (PT,4)
1860 PRINT D$;"READ CLASS-INFO,R";CREC;" ,B4"
1870 INPUT TCH(I)
1880 PRINT D$;"READ CLASS-INFO,R";CREC;" ,B9"
1890 INPUT RM(I)
1900 GOTO 1920
1910 TCH(I) = - 1:RM(I) = - 1
1920 NEXT I
1930 PRINT D$;"CLOSE CLASS-INFO"
1940 RETURN
1950 REM *****
1960 REM SUBROUTINE IN-ORDER
1970 REM *****
1980 PT = SP
1990 IF SNDX (PT,1) < > 0 THEN GOTO 2110
2000 LN = LN + 1
2010 ODR(LN) = PT
2020 D = 1
2030 GOSUB 2170
2040 D = 0
2050 PT = SNDX (PT,3)
2060 IF PT < > 0 THEN GOTO 1990
2070 GOSUB 2240
2080 IF D = 3 THEN GOTO 2160
2090 IF D = 1 THEN GOTO 2070
2100 GOTO 2000
2110 D = 2
2120 GOSUB 2170
2130 D = 0
2140 PT = SNDX (PT,1)
2150 GOTO 1990
2160 RETURN
2170 REM *****
2180 REM SUBROUTINE PUSH-STACK
2190 REM *****
2200 LST = LST + 1
2210 STACK (LST,1) = PT
2220 STACK (LST,2) = D
2230 RETURN
2240 REM *****
2250 REM SUBROUTINE PUSH-STACK
2260 REM *****
2270 IF LST = 0 THEN GOTO 2320
2280 PT = STACK (LST,1)
2290 D = STACK (LST,2)

```

```

2300 LST = LST - 1
2310 GOTO 2330
2320 D = 3
2330 RETURN
2335 REM
2340 REM *****
2350 REM SUBROUTINE PRINT MID-SEMESTER GRADES
2360 REM *****
2370 HOME : PRINT
2380 PRINT "MID-SEMSESTER GRADES FOR STUDENT: "; NUMB
2390 PRINT : PRINT
2400 PRINT "PER"; TAB( 8); "CLASS"; TAB( 16); "TEA"; TAB( 22);
2405 PRINT "RM"; TAB( 27); "GR"; TAB( 32); "COM"
2410 FOR I = 1 TO 7
2420 PRINT
2430 PRINT TAB( 1)I; TAB( 9)SCHD(I); TAB( 16)TCH(I);
2435 PRINT TAB( 22)RM(I); TAB( 27)GMT(I); TAB( 32)CMT(I)
2440 NEXT I
2450 PRINT : PRINT
2460 INPUT "PRESS RETURN TO CONTINUE: "; ZZ$
2470 RETURN
2480 REM *****
2490 REM SUBROUTINE TO PRINT FINIAL-GRADES
2500 REM *****
2510 HOME : PRINT
2520 PRINT "FINIAL GRADES FOR STUDENT: "; NUMB
2530 PRINT : PRINT
2540 PRINT TAB( 22)"MID"; TAB( 26)"MID"; TAB( 33)"FIN";
2545 PRINT TAB( 38)"FIN"
2550 PRINT "PER"; TAB( 6)"CLASS"; TAB( 13)"TEA"; TAB( 18)
2555 PRINT "RM"; TAB( 22)"GR"; TAB( 26)"COM"; TAB( 33)"GR";
2558 PRINT TAB( 38)"COM"
2560 FOR I = 1 TO 7
2570 PRINT
2580 PRINT TAB( 2)I; TAB( 7)SCHD(I); TAB( 13)TCH(I);
2585 PRINT TAB( 18)RM(I); TAB( 22)GMT(I); TAB( 27)CMT(I);
2587 PRINT TAB( 33)GF(I); TAB( 39)CF(I)
2590 NEXT I
2600 PRINT : PRINT
2610 INPUT "PRESS RETURN TO CONTINUE: "; ZZ$
2620 RETURN

```

APPENDIX H

PROGRAM RETRIEVE CLASS LISTS LISTING

```

100 REM PROGRAM TO RETREIVE CLASS LIST
110 REM
120 DIM CINDX(500,4)
130 DIM TINDX(50,4)
140 DIM INFO(5)
150 DIM LST(35)
160 DIM STACK(150,2)
170 DIM ODR(150)
180 D$ = CHR$(4)
190 GOSUB 440
200 GOSUB 540
210 HOME : PRINT : PRINT
220 PRINT "THIS PROGRAM WILL FIND CLASS LIST FOR:"
230 PRINT : PRINT
240 PRINT "1.....CLASS LIST (ONE CLASS)"
250 PRINT : PRINT
260 PRINT "2.....CLASS LISTS (ONE TEACHER)"
270 PRINT : PRINT
280 PRINT "3.....CLASS LISTS (ALL CLASSES)"
290 PRINT : PRINT
300 PRINT "4.....TO EXIT PROGRAM"
310 PRINT : PRINT
320 PRINT : PRINT
330 INPUT "THE NUMBER OF YOUR CHOICE";ANS
340 IF ANS = 1 THEN GOSUB 670
350 IF ANS = 2 THEN GOSUB 820
360 IF ANS = 3 THEN GOSUB 1160
370 IF ANS = 4 THEN GOTO 390
380 GOTO 210
390 REM EXIT PROGRAM AND RUN MENU
400 PRINT D$;"RUN MENU,D1"
410 REM *****
420 REM SUBROUTINE READ TEACHER INDEX
430 REM *****
440 PRINT D$;"OPEN TEACHER-INDEX,D2"
450 PRINT D$;"READ TEACHER-INDEX"
460 INPUT NTEA: INPUT SLP: INPUT KOBAL
470 IF NTEA = 0 THEN GOTO 520

```

```

480  FOR I = 1 TO NTEA
490  FOR J = 1 TO 4
500  INPUT TINDX(I,J)
510  NEXT J: NEXT I
520  PRINT D$;"CLOSE TEACHER-INDEX"
530  RETURN
540  REM *****
550  REM  SUBROUTINE TO READ CLASS-INDEX
560  REM *****
570  PRINT D$;"OPEN CLASS-INDEX"
580  PRINT D$;"READ CLASS-INDEX"
590  INPUT NCLASS: INPUT S2P: INPUT OCBAL
600  IF NCLASS = 0 THEN GOTO 650
610  FOR I = 1 TO NCLASS
620  FOR J = 1 TO 4
630  INPUT CINDX(I,J)
640  NEXT J: NEXT I
650  PRINT D$;"CLOSE CLASS-INDEX"
660  RETURN
670  REM *****
680  REM  SUBROUTINE FOR CLASS
690  REM *****
700  HOME : PRINT : PRINT
710  INPUT "CLASS NUMBER: ";NUMB
720  PT = S2P
730  GOSUB 1860
740  IF FLAG = 1 THEN GOTO 780
750  GOSUB 1470
760  GOSUB 1710
770  GOTO 810
780  PRINT : PRINT : PRINT "CLASS IS NOT IN THE DATA BASE"
790  FLAG = 0
800  PRINT:PRINT
850  INPUT "PRESS THE RETURN KEY TO CONTINUE:";ZZ$
810  RETURN
820  REM *****
830  REM  TEACHER SUBROUTINE
840  REM *****
850  HOME : PRINT : PRINT
860  PRINT "ENTER TEACHER NUMBER"
870  PRINT : PRINT : INPUT NUMB
880  PT = S1P
890  IF PT = 0 THEN GOTO 1110
900  IF NUMB < TINDX(PT,2) THEN IK = 1
910  IF NUMB > TINDX(PT,2) THEN IK = 3
920  IF NUMB = TINDX(PT,2) THEN 950

```



```

930 PT = TINDX(PT,IK)
940 GOTO 890
950 REC = TINDX(PT,4)
960 PRINT D$;"OPEN TEACHER-SCHEDULE,L31"
970 FOR I = 1 TO 7
980 BYT = 3 + (I - 1) * 4
990 PRINT D$;"READ TEACHER-SCHEDULE,R";REC;" ,B";BYT
1000 INPUT SCHD(I)
1010 NEXT I
1020 FOR JJ = 1 TO 7
1030 PT = S2P
1040 NUMB = SCHD(JJ)
1050 IF NUMB = 999 THEN GOTO 1090
1060 GOSUB 1860
1070 GOSUB 1470
1080 GOSUB 1710
1090 NEXT JJ
1100 GOTO 1150
1110 PRINT : PRINT : PRINT
1120 PRINT " TEACHER IS NOT IN THE DATA BASE"
1130 PRINT : PRINT
1140 INPUT "PRESS THE RETURN KEY TO CONTINUE:      ";ZZ$
1150 RETURN
1160 REM *****
1170 REM  PROCESS ALL TEACHERS-ALL CLASSES
1180 REM *****
1190 IF NTEA = 0 THEN GOTO 1390
1200 GOSUB 1990
1210 FOR IJ = 1 TO NTEA
1220 PRINT D$;"OPEN TEACHER-SCHEDULE,L31"
1230 FOR IK = 1 TO 7
1240 BYT = 3 + (IK - 1) * 4
1250 PRINT D$;"READ TEACHER-SCHEDULE,R";ODR(IJ);" ,B";BYT
1260 INPUT SCHD(IK)
1270 NEXT IK
1280 PT = S2P
1290 FOR J = 1 TO 7
1300 NUMB = SCHD(J)
1310 IF NUMB = 999 THEN GOTO 1350
1320 GOSUB 1860
1330 GOSUB 1470
1340 GOSUB 1710
1350 NEXT J
1360 NEXT IJ
1370 PRINT D$;"CLOSE TEACHER-SCHEDULE"
1380 GOTO 1450
1390 HOME : PRINT : PRINT : PRINT

```

```

1400 PRINT "THE TEACHER INDEX HAS NOT BEEN BIULD"
1410 PRINT : PRINT : PRINT
1420 PRINT "BUILD TEACHER-INDEX USING BUILD PROGRAM.
1430 PRINT : PRINT : PRINT
1440 INPUT "PRnESS RETURN TO CONTINUE ";ZZ$
1450 RETURN
1460 D$ = CHR$ (4)
1470 REM *****
1480 REM SUBROUTINE TO READ CLASS RECORD
1490 REM *****
1500 PRINT D$;"OPEN CLASS-INFO,L225"
1510 PRINT D$;"READ CLASS-INFO,R";REC;"",B0"
1520 INPUT INFO(1)
1530 PRINT D$;"READ CLASS-INFO,R";REC;"",B4"
1540 INPUT INFO(2)
1550 PRINT D$;"READ CLASS-INFO,R";REC;"",B7"
1560 INPUT INFO(3)
1570 PRINT D$;"READ CLASS-INFO,R";REC;"",B9"
1580 INPUT INFO(4)
1590 PRINT D$;"READ CLASS-INFO,R";REC;"",B12"
1600 INPUT INFO(5)
1610 IF INFO(5) = 0 THEN GOTO 1670
1620 FOR I = 1 TO INFO(5)
1630 BYT = 15 + (I - 1) * 6
1640 PRINT D$;"READ CLASS-INFO,R";REC;"",B";BYT
1650 INPUT LST(I)
1660 NEXT I
1670 PRINT D$;"CLOSE"
1680 GOTO 1700
1690 FLAG = 1
1700 RETURN
1710 REM *****
1720 REM DISPLAY SUBROUTINE
1730 REM *****
1740 HOME
1750 PRINT "CLASS"; TAB( 9)"TEACHER"; TAB( 19)"PERIOD";
1755 PRINT TAB( 28)"RM#"; TAB( 34)"CSIZE"
1760 PRINT : PRINT
1770 PRINT TAB( 1)INFO(1); TAB( 11)INFO(2);
1773 PRINT TAB( 21)INFO(3);
1775 PRINT TAB( 28)INFO(4); TAB( 35)INFO(5)
1780 PRINT : PRINT
1790 IF INFO(5) = 0 THEN GOTO 1840
1800 FOR I = 1 TO INFO(5)
1810 PRINT LST(I)
1820 NEXT I
1830 PRINT : PRINT

```

```

1840 INPUT "PRESS THE RETURN KEY TO CONTINUE:  ";ZZ$
1850 RETURN
1860 REM *****
1870 REM  SUBROUTINE TO FIND RECORD NUMBER"
1880 REM *****
1890 IF PT = 0 THEN GOTO 1950
1900 IF NUMB < CINDX(PT,2) THEN IK = 1
1910 IF NUMB > CINDX(PT,2) THEN IK = 3
1920 IF NUMB = CINDX(PT,2) THEN 1970
1930 PT = CINDX(PT,IK)
1940 GOTO 1890
1950 REM  NUMB IS NOT IN INDEX
1960 FLAG = 1
1970 REC = CINDX(PT,4)
1980 RETURN
1990 REM *****
2000 REM  SUBROUTINE IN ORDER
2010 REM *****
2020 LST = 0:PT = SLP
2030 IF TINDX(PT,1) < > 0 THEN GOTO 2190
2040 PRINT PT
2050 IY = IY + 1
2060 ODR(IY) = TINDX(PT,4)
2070 PRINT ODR(IY),IY
2080 D = 1
2090 REM  PUST THE STACK
2100 GOSUB 2270
2110 D = 0
2120 PT = TINDX(PT,3)
2130 IF PT < > 0 THEN GOTO 2030
2140 REM  POOP THE STACK
2150 GOSUB 2340
2160 IF D = 3 THEN GOTO 2240
2170 IF D = 1 THEN GOTO 2140
2180 GOTO 2050
2190 D = 2
2200 GOSUB 2270
2210 D = 0
2220 PT = TINDX(PT,1)
2230 GOTO 2030
2240 FOR KK = 1 TO 7
2250 PRINT ODR(KK);
2260 NEXT KK
2270 REM *****
2280 REM  SUBROUTINE PUSH
2290 REM *****
2300 LST = LST + 1

```

```
2310  STACK(LST,1) = PT
2320  STACK(LST,2) = D
2330  RETURN
2340  REM *****
2350  REM  SUBROUTINE POP
2360  REM *****
2370  IF LST = 0 THEN  GOTO 2420
2380  PT = STACK(LST,1)
2390  D = STACK(LST,2)
2400  LST = LST - 1
2410  GOTO 2430
2420  D = 3
2430  RETURN
```

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